SPEC SFS® 2014 SP2 User's Guide

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1 Quick Start Guide

The SPEC SFS 2014 benchmark is used to measure the maximum sustainable throughput that a storage solution can deliver. The benchmark is protocol independent. It will run over any version of NFS or SMB, clustered file systems, object oriented file systems, local file systems, or any other POSIX compatible file system. Because this tool runs at the application system call level, it is file system type agnostic. This provides strong portability across operating systems, and storage solutions. The SPEC SFS 2014 benchmark already runs on Linux, Windows 7, Windows 8, Windows 10, Windows Server 2008, Windows Server 2012, Windows Server 2016, Mac OS X, BSD, Solaris, and AIX, and can be used to test any of the file system types that these systems offer.

The SPEC SFS 2014 benchmark consists of multiple workloads which represent real data processing file system environments. Each of these workloads may be run independently. A SPEC SFS 2014 benchmark publication may use any one of these workloads. The SfsManager allows the user to input parameters, run the benchmark, and review the results for all workloads.

The benchmark runs on a group of workstations and measures the performance of the storage solution that is providing files to the application layer. Because the benchmark runs at the application system call level, all components of the storage solution impact the performance of the solution – this includes the load generators themselves as well as any physical or virtual hardware between the load generators and where the data ultimately rests on stable storage.





The minimal configuration consists of one load generating client, and one NFS, SMB, iSCSI, FCP, GPFS, or Luster server. For Windows-based environments, a separate prime client is required for a minimal configuration.

The steps to produce a SPEC SFS 2014 result are:

- Install the SPEC SFS 2014 benchmark on the load generators
- Edit the configuration file on the prime client
- Configure the solution for testing
- Start the benchmark

1.1 Prerequisites

- Python version 2.7 or later must be installed on the Prime Client system.
- Matplotlib must be installed on the Prime Client system.
- See: <u>http://www.Matplotlib.org</u>
- Matplotlib dependencies (dateutil, numpy, pyparsing & six) must also be installed. dateutil pip install python-dateutil Numpy <u>http://www.Numpy.org</u> Pyparsing ... easy_install pyparsing Six <u>https://pypi.python.org/pypi/six</u>
- The test file systems must have the permissions set correctly in order to allow access by the clients.
- The test file systems must be mounted or mapped prior to execution of the benchmark. For Windows shares one may use UNC paths without mapping the shares.
- There must be interconnect/network connectivity between any storage system and clients, and between the clients and the Prime Client. The Prime Client is simply the system on which the benchmark run is started, and could be one of the clients. The prime client may also present load, in some configurations. When running on Windows, the prime client cannot also generate load, so at least one additional client is required.
- The contents of the SPEC SFS 2014 benchmark distribution must be accessible on all the systems where the benchmark will be installed.
- Using these quick start procedures assumes that the pre-compiled C code binaries, shipped with the benchmark, will be used.

1.2 Installing the SPEC SFS 2014 Benchmark

The SPEC SFS 2014 benchmark can be installed on client machines running either a UNIX-based or Windows operating system. Each of these require slightly different configuration and are described separately below.

UNIX client installation and configuration:

- Ensure that DNS is correctly configured.
 - Ensure that the client's hostname does not appear in the hosts file on either the 127.0.0.1 or ::1 line!
- Ensure that any/all forms of firewalls are disabled on all of the clients. (SELinux, iptables...)
- Ensure that all clients have ssh installed and configured such that clients can start commands on each other without any password challenges. (Example: ssh hostname command) Setup the ssh keys, and permit empty passwords.
- Install Python 2.7, and ensure that python is in the user's search path. Python may be downloaded from http://www.python.org.
- Install the SPEC SFS 2014 benchmark using the following steps:
 - Login to the client (as root)
 - o Download or copy the SPEC SFS 2014 ISO or archive to the client
 - o Loop mount the ISO or expand the archive
 - cd to the top level directory containing the SPEC SFS 2014 contents
 - Enter 'python SfsManager --install-dir="*destination_directory*"' (where "*destination_directory*", without the double-quotes, is the full path where you wish to have the benchmark installed)

• It is strongly recommended to apply the recommended tuning for somaxconn in 11.1 below for Linux hosts

Windows client installation and configuration:

- Ensure that DNS is correctly configured.
 - Ensure that the client's hostname does not appear in the hosts file on either the 127.0.0.1 or ::1 line!
- Ensure that all Windows clients are properly configured in the Active Directory Domain.
- Install Python 2.7 or later, and ensure that python is in the user's search path. Python may be downloaded from http://www.python.org
- Install the SPEC SFS 2014 benchmark
 - Start a command prompt window. This can be done using the 'Start' button, choosing 'Run...' and entering 'cmd'.
 - \circ $\,$ Download or copy the SPEC SFS 2014 ISO or archive to the client
 - Attach the ISO as a virtual optical drive or expand the archive
 - o chdir to the top level directory containing the SPEC SFS 2014 directory
 - Enter 'python SfsManager --install-dir='"*destination_directory*" (where "*destination_directory*", enclosed by double-quotes, is where you wish to have the benchmark installed)

Note: Please use a destination_directory that is *not* in the user's home directory. The actual path to a User's home directory location varies widely across versions of Windows and can make heterogeneous clients problematic.

1.3 Editing the configuration file on the prime client

On the prime client, the default configuration file is called sfs_rc. The user must edit the sfs_rc configuration file, and only needs to edit it on the prime client. The user does not need to edit, or even have, a configuration file on the other load generating clients. On the prime client, edit the values for:

• BENCHMARK

Name of the benchmark to run. Valid values are: **DATABASE**, **EDA**, **SWBUILD**, **VDA**, or **VDI**.

• LOAD

Each workload has an associated business metric as a unit of workload. The magnitude of the workload to run is specified with the **LOAD** parameter in units of the workload's business metric. Valid values for **LOAD** are either a starting number or a list of values, all positive integers. If a single value is specified, it is interpreted as a starting value and used in conjunction with **INCR_LOAD** and **NUM_RUNS**. The following table shows the name for the business metric corresponding to each workload type.

Workload	Business Metric (LOAD parameter)
DATABASE	DATABASES
EDA	JOB_SETS
SWBUILD	BUILDS
VDA	STREAMS
VDI	DESKTOPS

If a list of values if specified, at least 10 uniformly spaced data points must be specified for valid benchmark execution. For more detail on the requirements for uniformly spaced data points, see section 5.3 "Data Point Specification for Results Disclosure" in the SPEC SFS® 2014 SP2 Run and Reporting Rules.

As a helpful guideline, some rules of thumb for the resources required per business metric for the different workloads:

Capacity requirements per business metric:				
= 24 GiB per DATABASE				
= 11 GiB per JOB_SET				
= 5 GiB per BUILD				
= 24 GiB per STREAM				
= 12 GiB per DESKTOP				

Client memory requirements per business metric: DATABASE = 55 MiB per DATABASE

EDA	= 12 MiB per JOB_SET
SWBUILD	= 400 MiB per BUILD
VDA	= 10 MiB per STREAM
VDI	= 8 MiB per DESKTOP

• INCR_LOAD

Incremental increase in load for successive data points in a run. This parameter is used only if **LOAD** consists of a single (initial) value. To ensure equally spaced points, the value of **LOAD** and **INCR_LOAD** must be equal.

• NUM_RUNS

The number of load points to run and measure (minimum of 10 for a publishable result). This parameter is used only if **INCR_LOAD** is specified.

• CLIENT_MOUNTPOINTS

The list of local mount points, local directories, or shares, to use in the testing. The value of CLIENT_MOUNTPOINTS can take several different forms:

- UNIX style: client:/exportfs1 client:/exportfs2 ...
 - Used for local storage or mounted network shares
- Windows style: client:\\server\exportfs1 client:\\server\exportfs2 ...
- Use a file that contains the mount points: mountpoints_file.txt
 - For more detail, see section 3.3 "*Configuring the Required Benchmark Parameters*" below
- The CLIENT_MOUNTPOINTS list is used to assign both a load generator as well as a storage location to a single business metric. The order of entries in the CLIENT_MOUNTPOINTS list matters, as the list is consumed sequentially to assign a client and a storage location to each business metric as the load scales up. Once the list of CLIENT_MOUNTPOINTS is exhausted, the list will be reused from the start as many times as necessary.
- The order of the CLIENT_MOUNTPOINTS list will affect how the load scales up in the SUT be sure to order this list to avoid artificial bottlenecks. For example, if the SUT has many load generators and many file systems accessible by all load generators, it may be desirable to order the CLIENT_MOUNTPOINTS list so that a unique load generator and unique file system is used with each CLIENT_MOUNTPOINTS entry until re-use is required to exhaust all possible combinations. By spreading the load as evenly as possible, it is possible to avoid one load generator or one file system becoming a hot-spot.
- The ideal way to order the CLIENT_MOUNTPOINTS list for a given SUT will depend heavily on its architecture.

• EXEC_PATH

The full path to the SPEC SFS 2014 executable. Currently the executable is called *netmist* for POSIX systems and *netmist.exe* for Windows systems.

• USER

The user account name, which must be configured on all clients, to be used for the benchmark execution. To specify a domain, prefix the user with the domain name, separated by a backslash. E.g. DOMAIN\User33

• WARMUP_TIME

The amount of time, in seconds, that the benchmark will spend in WARMUP before initiating the measurement ("RUN") phase. The minimum for publishable submissions is 300 seconds (five minutes). The maximum value for a publishable submission is 604,800 seconds (one week). If this value is modified, the value used must be noted in the Other Solution Notes (otherSutNotes) field of the disclosure.

• IPV6_ENABLE

Set to "1" or "Yes" when the benchmark should use IPv6 to communicate with other benchmark processes.

• **PRIME_MON_SCRIPT**

The name of a shell script or other executable program which will be invoked to control any external programs. These external programs must be performance neutral and their actions must comply with the SPEC SFS® 2014 SP2 Run and Reporting Rules. If this option is used, the executable or script used must be disclosed. If this option is used, the script must be noted in the Other Solution Notes (otherSutNotes) field of the disclosure. Scripts may also be attached as a config diagram must be referenced from the Other Solution Notes (otherSutNotes) field.

An example of how to write a script to be invoked by PRIME_MON_SCRIPT, see the sfs_ext_mon example script included with the benchmark.

• PRIME_MON_ARGS

Arguments which are passed to the executable specified in PRIME_MON_SCRIPT. If this option is used, the values used must be noted in the Other Solution Notes (otherSutNotes) field of the disclosure. Each argument passed via PRIME_MON_ARGS appears in a separate command line argument to PRIME_MON_SCRIPT – there is no escaping, etc. to encapsulate all PRIME_MON_ARGS into one argument to PRIME_MON_SCRIPT.

• NETMIST_LOGS

Set the path to the directory in which to store log files from the load generators. The same path will be used on all clients. If this path is not set, /tmp/ or c:\tmp\ will be used.

• PASSWORD

The password for the user specified in USER. (Only applicable when running on Windows platforms)

1.4 Installing the license key

• Obtain your license number from SPEC. Create the *netmist_license_key* file in either /tmp or in the current working directory where you will run the benchmark. This file should be a simple text file that contains: LICENSE KEY ###### Where the ##### is the license number that you have received from SPEC.

You must create this file before running the benchmark.

1.5 Configuring the storage solution for testing

- Mount all working directories on the clients (POSIX only). The path names must match the values specified in the CLIENT_MOUNTPOINTS parameter in the SPEC SFS 2014 configuration file. Mapping the shares is not needed if running on Windows and using UNC paths.
- Ensure the exported file systems have read/write permissions.
- Ensure access is permitted for username, password, and domain. (SMB testing only)

1.6 Starting the benchmark

Note that the SfsManager must be run under the same user id (UID) on the all of the clients, including the prime client.

- Change directories to the *destination_directory* specified during install.
- On the Prime client:
 - Enter 'python SfsManager -r sfs_config_file -s output_files_suffix

1.7 Monitoring the benchmark execution

• On the Prime client, change directories to the *destination_directory* from the installation step above, then 'cd result'

The user may now examine the benchmark logs, as well as the results. As the benchmark runs, the results are stored in the files with names like:

sfssum_*.{txt,xml}	Summary file used in the submission process described later.
sfslog_*.log	Log file of the current activity.

After all load points are complete, the results from each client are collected into the result directory on prime client. The client logs are files with names like: sfsc001.*

sfsc*.* The client log files.

More detailed client logs can be found on each client in the path specified by the NETMIST_LOGS RC file parameter, or /tmp/ or C:\tmp\ by default. It is recommended that these log files be purged from each client between each run of the benchmark – you may wish to save these with the other log files from the run before deleting them.

1.8 Examining the results after the benchmark execution has completed

The results of the benchmark are summarized in the sfssum.* files in the result directory on the prime client. These may be examined with any text editing software package. The sfssum_*.xml is the XML summary file that will be used for the submission process, described later in this document.

2 Introduction

The SPEC SFS 2014 benchmark is the latest version of the Standard Performance Evaluation Corporation benchmark that measures storage solution throughput and response time. It provides a standardized method for comparing performance across different vendor platforms.

This document specifies how the SPEC SFS 2014 benchmark is to be run for measuring and publicly reporting performance results, and includes a guide to using the SFS tools. The SPEC SFS® 2014 SP2 Run and Reporting Rules (included in a separate companion document that is included in the SPEC SFS 2014 distribution) have been established by the SPEC Storage Subcommittee and approved by the SPEC Open Systems Steering Committee. They ensure that results generated with this suite are meaningful, comparable to other generated results, and are repeatable. Per the SPEC license agreement, all results publicly disclosed must adhere to these Run and Reporting Rules.

SPEC requires that any public use of results from this benchmark follow the <u>SPEC OSG Fair Use Policy</u> at <u>https://www.spec.org/fairuse.html</u>. In the case where it appears that these guidelines have not been adhered to, SPEC may investigate and request that the published material be corrected.

The initial SPEC SFS 2014 release of the benchmark included major workload and functionality changes, as well as clarification of run rules. The code changes were NOT performance neutral, therefore comparing SPEC SFS 2014 results with results from previous major versions of the SFS benchmark (e.g. SFS 2008, SFS 97_R1) is NOT allowed.

The SPEC SFS 2014 SP2 benchmark release includes existing workloads and one new workload. For the existing workloads (DATABASE, SWBUILD, VDA, VDI), the code changes between SP2 and SP1 are performance-neutral. This release also adds the new EDA workload.

2.1 What is the SPEC SFS 2014 Benchmark

The SPEC SFS 2014 benchmark is used to measure the maximum sustainable throughput that a storage solution can deliver. The benchmark is protocol independent. It will run over any version of NFS or SMB, clustered file systems, object oriented file systems, local file systems, or any other POSIX compatible file system. Because this tool runs at the application system call level, it is file system type agnostic. This provides strong portability across operating systems, and storage solutions. The SPEC SFS 2014 benchmark already runs on Linux, Windows Vista, Windows 7, Windows 8, Windows 10, Windows Server 2008, Windows Server 2012, Windows Server 2016, Mac OS X, BSD, Solaris, and AIX, and can be used to test any of the file system types that these systems offer.

The SPEC SFS 2014 workloads are a mixture of file meta-data and data oriented operations. The SPEC SFS 2014 benchmark is fully multi-client aware, and is a distributed application that coordinates and conducts the testing across all of the client nodes that are used to test a storage solution.

The benchmark runs on a group of workstations and measures the performance of the storage solution that is providing files to the application layer on the workstations. Each workload consists of several typical file operations. The following is the current set of operations that are measured.

read()	Read file data sequentially
read_file()	Read an entire file sequentially
mmap_read()	Read file data using the mmap() API.
read_random()	Read file data at random offsets in the files
write()	Write file data sequentially
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write_file()	Write an entire file sequentially.
mmap_write()	Write a file using the mmap() API.
write_random()	Write file data at random offsets in the files.
rmw()	Read+modify+write file data at random offsets in files.
mkdir()	Create a directory
rmdir()	Removes a directory
unlink()	Unlink/remove an empty file.
unlink2()	Unlink/remove a non-empty file.
append()	Append to the end of an existing file.
lock()	Lock a file.
unlock()	Unlock a file.
access()	Perform the access() system call on a file.
stat()	Perform the stat() system call on a file.
chmod()	Perform the chmod() system call on a file.
create()	Create a new file.
readdir()	Perform a readdir() system call on a directory.
statfs()	Perform the statfs() system call on a filesystem.
copyfile()	Copy a file.
rename()	Rename a file
pathconf()	Perform the pathconf() system call.

The read() and write() operations are performing sequential I/O to the data files. The read_random() and write_random() perform I/O at random offsets within the files. The read_file and write_file operate in whole files. Rmw is a read_modify_write operation.

The results of the benchmark are:

- 1. Maximum workload-specific Business Metric achieved.
- 2. Aggregate Ops/sec that the storage solution can sustain at requested or peak load.
- 3. Average file operation latency in milliseconds.
- 4. Aggregate KiB/sec that the storage solution can sustain at requested or peak load.

The SPEC SFS 2014 benchmark includes multiple workloads:

DATABASE	Transactional SQL database
EDA	Electronic design automation
SWBUILD	Software build
VDA	Video data acquisition
VDI	Virtual desktop infrastructure

The user has the option to submit results using any or all of the above workloads. The SfsManager accepts benchmark run configuration information, starts benchmark execution, and collects results from the SPEC SFS 2014 benchmark.

2.2 SPEC SFS 2014 Benchmark Overview

The SPEC SFS 2014 benchmark is used to test the performance capability of storage solutions. Performance is measured in metrics related to the specific workload tested.

In a typical SPEC SFS test configuration, a series of load generating clients are directed through a network at file systems shared or exported from a file server. The SUT consists of the load generators, the network, the file server, and finally the stable backend storage. The SUT is the entire solution under test. More details of the SUT may be found in section 6.4 of the SPEC SFS® 2014 SP2 Run and Reporting Rules document.

2.2.1 Benchmark Manager

2.2.1.1 Introduction

The benchmark manager is called the SfsManager. It is a Python program that requires Python version 2.7 or higher. It is recommended to get Python from http://www.python.org/

You can get the syntax by running "python SfsManager -h"

Usage: python SfsManager [options]

Command line option: Required for Benchmark Execution: [-r <file>] or [--rc-file=<file>]...... Specify rc file Required for Benchmark Installation: Optional: (default=sfs2014_SP2) [-b <file>] or [--benchmark-file=<file>]..... benchmark definition file [-d <dir>] or [--results-dir=<dir>]..... Results directory, use full path [-i] or [--ignore-override]...... Bypass override of official workload parameters [-e <filename>] or [--export=<filename>]..... Export workload definitions to a file [-a]..... Auto mode (max): finds maximum passing load value [-A <scaling>]..... Auto mode (curve): generate 10 point curve based on result of -a (if used) or LOAD [-A] (continued)......<scaling> is a percentage between 0-100 to scale the maximum LOAD. [-A] (continued)...... If -A and -a are used together a '<suffix>.auto' is used for finding the maximum. [--save-config=<file>]...... Save the token config file and executeable command line args [--test-only]...... Simulate a set of load points without actually running. [-h]..... Show usage info [-v]...... Show version number [--debug]..... Detailed output

The only required parameter is an rc file. The -a and -A options to SfsManager are not valid for publication runs. These options are included as a convenience only. The base directory has an example called sfs_rc, whose contents are

#
#
sfs_rc
#
Specify netmist parameters for generic runs in this file.
#
The following parameters are configurable within the SFS run and
reporting rules.
#
Official BENCHMARK values are
-SWBUILD
-VDA
-VDI

-DATABASE # -EDA # BENCHMARK=VDI LOAD=10 INCR_LOAD=10 NUM RUNS=10 CLIENT MOUNTPOINTS= EXEC PATH=/usr/local/bin/netmist USER=root WARMUP TIME=300 IPV6 ENABLE=0 PRIME MON SCRIPT= PRIME MON ARGS= NETMIST LOGS= INIT RATE=0 # # Specifying a password is only required for Windows clients # PASSWORD= # # DO NOT EDIT BELOW THIS LINE FOR AN OFFICIAL BENCHMARK SUBMISSION # # Constraints and overrides on the values below this line can be found in the # benchmark XML file (default is benchmarks.xml). To bypass all overrides # use the --ignore-overrides flag in SfsManager. Using the flag will make # the results invalid for formal submission. # RUNTIME=300 WORKLOAD FILE= **OPRATE MULTIPLIER=** CLIENT_MEM=1g AGGR_CAP=1g FILE SIZE= DIR_COUNT=10 FILES PER DIR=100 UNLINK FILES=0 LATENCY_GRAPH=1 HEARTBEAT_NOTIFICATIONS=1 DISABLE FSYNCS=0 USE RSHRCP=0 BYTE OFFSET=0 MAX FD= PIT_SERVER= PIT PORT= LOCAL_ONLY=0 FILE_ACCESS_LIST=0 SHARING MODE=0 SOCK_DEBUG=0 13 SPEC SFS® 2014 SP2

TRACEDEBUG=0 NO_OP_VALIDATE=0 NO_SHARED_BUCKETS=0 UNLINK2_NO_RECREATE=0

Some things to keep in mind:

- You never need single or double quotes around anything
- The only parameters that must be explicitly specified are BENCHMARK, LOAD, CLIENT_MOUNTPOINTS, USER, EXEC_PATH, and PASSWORD (Windows only)
- The binary parameters all get translated to 1 or 0, but you can also use "yes", "no", "Y", "N", "on", or "off"

2.2.1.2 The Summary Files

Each run will produce 2 summary files: sfssum_<suffix>.txt and sfssum_<suffix>.xml. The txt file is a human readable summary file with the following fields

Field	Value
1	Business metric (may be blank for custom workloads)
2	Requested op rate (blank if not set)
3	Achieved op rate in Ops/s
4	Average latency in milliseconds
5	Total throughput in KiB/s
6	Read Throughput in KiB/s
7	Write Throughput in KiB/s
8	Run time in seconds
9	# of clients
10	# of procs per client
11	Average file size in KiB
12	Client data set size in MiB
13	Total starting data set size in MiB
14	Total initial file set space in MiB
15	Max file space in MiB
16	Workload name
17	Run validity field (blank if valid or no validation criteria exist in custom workload)

2.2.1.3 Example 1: Official Benchmark Run - VDI

In this example there are 2 clients, client1 and client2, each with a single mountpoint /mnt. The user starts with 50 DESKTOPS and generates 10 load points while incrementing the number of DESKTOPS by 50 for each load point. The RC file, vdi_rc, contains the following

Contents of vdi_rc

BENCHMARK=VDI LOAD=50 INCR_LOAD=50 NUM_RUNS=10 CLIENT_MOUNTPOINTS=client1:/mnt client2:/mnt EXEC_PATH=/usr/local/bin/netmist USERNAME=sfsuser

From the prime client, which may or may not be one of the 2 clients listed in CLIENT_MOUNTPOINTS, the user starts the benchmark with:

[prime client]\$ python SfsManager -r vdi_rc -s vditest

After the benchmark completes, the results directory will contain the following files:

File	Description
sfslog_vditest.log	Overall log file
sfssum_vditest.txt	Summary file in text format
sfssum_vditest.xml	Summary file in XML format
sfsc001.vditest	Detailed results for client1
sfsc002.vditest	Detailed results for client2

2.2.2 Database (DATABASE) Benchmark

2.2.2.1 Workload description

This workload represents the typical behavior of a transactional SQL database. The complete workload is a mixture of DB_TABLE and DB_LOG workloads. The DB_TABLE workload is the database component, and DB_LOG represents the log writer component of a database operation.

2.2.2.2 Workload characteristics

2.2.2.1 DATABASE File Operation Distribution

	Operation	%	Operation	%
	read	1	read file	0
Ę	mmap read	0	rand read	79
utic	write	0	write file	0
ribu	mmap write	0	rand write	20
Dist	rmw	0	append	0
u C	mkdir	0	rmdir	0
atio	readdir	0	create	0
per	unlink	0	unlink2	0
e O	stat	0	access	0
ï	rename	0	copyfile	0
	locking	0	chmod	0
	statfs	0	pathconf	0

2.2.2.1.1 DB_TABLE File Operation Distribution

2.2.2.1.2 DB_LOG File Operation Distribution

	Operation	%	Operation	%
	read	0	read file	0
u o	mmap read	0	rand read	0
uti	write	80	write file	0
trib	mmap write	0	rand write	20
Dis	rmw	0	append	0
on	mkdir	0	rmdir	0
rati	readdir	0	create	0
be	unlink	0	unlink2	0
e O	stat	0	access	0
ii.	rename	0	copyfile	0
	locking	0	chmod	0
	statfs	0	pathconf	0

2.2.2.2 DATABASE Read Transfer Size Distribution

	Slot	Start	End	%
	0	1	511	0
	1	512	1023	0
_	2	1024	2047	0
tio	3	2048	4095	0
in qi	4	4096	4096	0
istri	5	4097	8191	0
ē	6	8192	8192	99
Siz	7	8193	16383	0
sfer	8	16384	16384	0
lans	9	16385	32767	0
F P	10	32768	32768	0
Rea	11	65536	65536	0
	12	98304	98304	0
	13	131072	131072	0
	14	262144	262144	0
	15	1048576	1048576	1

2.2.2.2.1 DB_TABLE Read Transfer Size Distribution

2.2.2.2.2 DB_LOG Read Transfer Size Distribution

	Slot	Start	End	%
	0	1	511	0
	1	512	1023	0
c	2	1024	2047	0
Itio	3	2048	4095	0
ibu	4	4096	4096	0
isti	5	4097	8191	0
e e	6	8192	8192	99
Siz	7	8193	16383	0
sfer	8	16384	16384	0
aus	9	16385	32767	0
Ē,	10	32768	32768	0
(eac	11	65536	65536	0
<u> </u>	12	98304	98304	0
	13	131072	131072	0
	14	262144	262144	0
	15	1048576	1048576	1

2.2.2.3 DATABASE Write Transfer Size Distribution

	Slot	Start	End	%
	0	512	512	0
	1	1024	1024	0
_	2	1536	1536	0
tion	3	2048	2048	0
ibu	4	2560	2560	0
istr	5	3072	3072	0
e D	6	3584	3584	0
r Siz	7	4096	4096	0
sfei	8	4608	4608	0
ran	9	5120	5120	0
e T	10	8192	8192	100
Vrit	11	12288	12288	0
>	12	16384	16384	0
	13	20480	20480	0
	14	24576	24576	0
	15	32768	32768	0

2.2.2.3.1 DB_TABLE Write Transfer Size Distribution

2.2.2.3.2 DB_LOG Write Transfer Size Distribution

	Slot	Start	End	%
	0	512	512	5
	1	1024	1024	5
Ę	2	1536	1536	5
Itio	3	2048	2048	5
ribu	4	2560	2560	5
)ist	5	3072	3072	5
e L	6	3584	3584	5
Siz	7	4096	4096	5
sfei	8	4608	4608	5
ran	9	5120	5120	5
еŢ	10	8192	8192	10
Vrit	11	12288	12288	10
5	12	16384	16384	10
	13	20480	20480	10
	14	24576	24576	10
	15	32768	32768	0

2.2.2.4 DATABASE Miscellaneous Characteristics

	Option	Value	Option	Value
ns	write commit %	100	background	0
Dec	% direct	100	sharemode	1
llar	% osync	0	uniform size dist	1
sce	% notification	0	init rate throttle	0
Σ	LRU	0	init read flag	1
	release version	1		

2.2.2.4.2 DB_LOG Miscellaneous Characteristics

	Option	Value	Option	Value
sno	write commit %	100	background	0
nec	% direct	100	sharemode	1
ella	% osync	0	uniform size dist	1
isce	% notification	0	init rate throttle	0
Σ	LRU	0	init read flag	1
	release version	1		

2.2.2.5 DATABASE Access Pattern Characteristics

2.2.2.5.1 DB_TABLE Access Pattern Characteristics

rns	Option	Value	Option	Value
ttei	rand dist behavior	2	% per spot	5
ba Ba	min acc per spot	1000	access mult spot	5
ces	affinity %	20	spot shape	0
Ac	geometric %	2	align	0

2.2.2.5.2 DB_LOG Access Pattern Characteristics

		Option	Value	Option	Value
SS	rns	rand dist behavior	2	% per spot	5
CC G	ttei	min acc per spot	1000	access mult spot	5
Pat Pat	Pai	affinity %	20	spot shape	0
		geometric %	2	align	0

2.2.2.6 DATABASE Content Pattern Characteristics

2.2.2.6.1 DB_TABLE Access Pattern Characteristics

S	Option	Value	Option	Value
ter	dedup %	0	dedup within %	0
Pat	dedup across %	0	dedup group count	1
ent	dedup granule size	4096	dedup gran rep limit	100
nte	compress %	50	comp granule size	8192
ပိ	cipher flag	0	pattern version	1

2.2.2.6.2 DB_LOG Access Pattern Characteristics

ns	Option	Value	Option	Value
ter	dedup %	0	dedup within %	0
Pat	dedup across %	0	dedup group count	1
ent	dedup granule size	4096	dedup gran rep limit	100
onte	compress %	50	comp granule size	8192
č	cipher flag	0	pattern version	1

2.2.2.7 DATABASE Execution Parameters

2.2.2.7.1 DB_TABLE Execution Parameters

ers	Parameter	Value	Parameter		Value
utic nete	Procs	10	Dirs per proc		1
ram	Oprate per proc	16	Files per dir		5
a E	Avg file size			20	0 MiB

2.2.2.7.2 DB_LOG Execution Parameters

ers	Parameter	Value	Parameter		Value
utio	Procs	1	Dirs per proc		1
feci	Oprate per proc	32	Files per dir		5
Pa E	Avg file size			20	0 MiB

2.2.2.8 DATABASE Validation Criteria

Threshold	Value
Per proc oprate	>= 75% of requested
Global average oprate	>= 95% of requested
Per proc maximum acceptable latency	N/A
Global average maximum acceptable latency	N/A
Workload oprate ratio variance	<= +/- 5%

2.2.3 Electronic Design Automation (EDA) Benchmark

2.2.3.1 Workload description

This workload represents the typical behavior of a mixture of EDA applications. The complete workload is a mixture of EDA_FRONTEND and EDA_BACKEND workloads. The EDA_FRONTEND workload is the EDA frontend processing applications, and EDA_BACKEND represents the EDA backend applications that generate the final output files.

2.2.3.2 Workload characteristics

2.2.3.2.1 EDA File Operation Distribution

	Operation	%	Operation	%
	read	0	read file	7
u n	mmap read	0	rand read	8
rtic	write	0	write file	10
ribı	mmap write	0	rand write	15
Dist	rmw	0	append	0
] uc	mkdir	1	rmdir	0
atio	readdir	0	create	2
ber	unlink	1	unlink2	1
e O	stat	39	access	15
Ξ	rename	0	copyfile	0
	locking	0	chmod	1
	statfs	0	pathconf	0

2.2.3.2.1.1 EDA_FRONTEND File Operation Distribution

2.2.3.2.1.2 EDA_BACKEND File Operation Distribution

	Operation	%	Operation	%
	read	50	read file	0
u o	mmap read	0	rand read	0
uti	write	50	write file	0
trib	mmap write	0	rand write	0
Dis	rmw	0	append	0
on	mkdir	0	rmdir	0
rati	readdir	0	create	0
be	unlink	0	unlink2	0
e O	stat	0	access	0
Ē	rename	0	copyfile	0
	locking	0	chmod	0
	statfs	0	pathconf	0

2.2.3.2.2 EDA Read Transfer Size Distribution

	Slot	Start	End	%
	0	1	511	4
	1	512	1023	0
_	2	1024	2047	0
tio	3	2048	4095	2
in qi	4	4096	8191	43
istri	5	8192	16383	30
۵ آ	6	16384	32767	21
·Siz	7	32768	65535	0
sfer	8	65536	65536	0
lan	9	131072	131072	0
F P	10	1	1	0
Rea	11	1	1	0
	12	1	1	0
	13	1	1	0
	14	1	1	0
	15	1	1	0

2.2.3.2.2.1 EDA_FRONTEND Read Transfer Size Distribution

2.2.3.2.2.2 EDA_BACKEND Read Transfer Size Distribution

	Slot	Start	End	%
	0	1	511	0
	1	512	1023	0
c	2	1024	2047	0
Itio	3	2048	4095	0
ribu	4	4096	8191	0
ist	5	8192	16383	0
e e	6	16384	32767	0
· Siz	7	32768	65535	49
sfer	8	65536	65536	51
ans	9	131072	131072	0
Ē R	10	1	1	0
(eac	11	1	1	0
<u></u>	12	1	1	0
	13	1	1	0
	14	1	1	0
	15	1	1	0

2.2.3.2.3 EDA Write Transfer Size Distribution

	Slot	Start	End	%
	0	1	511	25
	1	512	1023	10
-	2	1024	2047	15
tio	3	2048	4095	18
ibu	4	4096	8191	27
istr	5	8292	16383	3
e D	6	16384	32767	2
r Siz	7	32768	65535	0
sfei	8	65536	65536	0
ran	9	131072	131072	0
e T	10	1	1	0
Vrit	11	1	1	0
>	12	1	1	0
	13	1	1	0
	14	1	1	0
	15	1	1	0

2.2.3.2.3.1 EDA_FRONTEND Write Transfer Size Distribution

2.2.3.2.3.2 EDA_BACKEND Write Transfer Size Distribution

	Slot	Start	End	%
	0	1	511	0
	1	512	1023	0
Ę	2	1024	2047	0
ltio	3	2048	4095	0
ribu	4	4096	8191	0
Jist	5	8292	16383	0
Se 🛛	6	16384	32767	0
r Si	7	32768	65535	45
sfei	8	65536	131072	55
ran	9	131072	131072	0
еT	10	1	1	0
Vrit	11	1	1	0
>	12	1	1	0
	13	1	1	0
	14	1	1	0
	15	1	1	0

2.2.3.2.4 EDA Miscellaneous Characteristics

	Option	Value	Option	Value
ns	write commit %	15	background	0
Jeo J	% direct	0	sharemode	0
llar	% osync	0	uniform size dist	0
sce	% notification	0	init rate throttle	1
Ξ	LRU	1	init read flag	0
	release version	2		

2.2.3.2.4.1 EDA_FRONTEND Miscellaneous Characteristics

2.2.3.2.4.2 EDA_BACKEND Miscellaneous Characcteristics

	Option	Value	Option	Value
SUC	write commit %	15	background	0
neo	% direct	50	sharemode	0
allai	% osync	5	uniform size dist	0
isce	% notification	0	init rate throttle	1
Ξ	LRU	1	init read flag	0
	release version	2		

2.2.3.2.5 EDA Access Pattern Characteristics

2.2.3.2.5.1 EDA_FRONTEND Access Pattern Characteristics

ns	Option	Value	Option	Value
ttei	rand dist behavior	0	% per spot	0
, Pa	min acc per spot	0	access mult spot	5
Ces	affinity %	0	spot shape	0
Ac	geometric %	50	align	0

2.2.3.2.5.2 EDA_BACKEND Access Pattern Characteristics

		Option	Value	Option	Value
Access Patterns	rns	rand dist behavior	0	% per spot	0
	ttei	min acc per spot	0	access mult spot	5
	Pa	affinity %	0	spot shape	0
		geometric %	50	align	0

2.2.3.2.6 EDA Content Pattern Characteristics

S	Option	Value	Option	Value
ter	dedup %	50	dedup within %	0
Pat	dedup across %	0	dedup group count	1
ent	dedup granule size	4096	dedup gran rep limit	100
nte	compress %	50	comp granule size	8192
ပိ	cipher flag	0	pattern version	2

2.2.3.2.6.1 EDA_FRONTEND Content Pattern Characteristics

2.2.3.2.6.2 EDA_BACKEND Content Pattern Characteristics

ns	Option	Value	Option	Value
ter	dedup %	40	dedup within %	0
Pat	dedup across %	0	dedup group count	1
ent	dedup granule size	4096	dedup gran rep limit	100
nte	compress %	20	comp granule size	8192
ö	cipher flag	0	pattern version	2

2.2.3.2.7 EDA Execution Parameters

2.2.3.2.7.1 EDA_FRONTEND Execution Parameters

Execution Parameters	ers	Parameter	Value	Parameter	Value
	lete	Procs	3	Dirs per proc	10
	ran	Oprate per proc	100	Files per dir	10
	Ра	Avg file size			16 KiB

2.2.3.2.7.2 EDA_BACKEND Execution Parameters

2	ราร	Parameter	Value	Parameter		Value
Executio Paramete	Jet	Procs	2	Dirs per proc		5
	ran	Oprate per proc	75	Files per dir		10
	Avg file size			1	0 MiB	

2.2.3.2.8 EDA Validation Criteria

Threshold	Value
Per proc oprate	>= 75% of requested
Global average oprate	>= 95% of requested
Per proc maximum acceptable latency	N/A
Global average maximum acceptable latency	N/A
Workload oprate ratio variance	<= +/- 5%

2.2.4 Software Build (SWBUILD) Benchmark

2.2.4.1 Benchmark Description

The software build type workload is a classic meta-data intensive build workload. This workload was derived from analysis of software builds, and traces collected on systems in the software build arena. Conceptually, these tests are similar to running unix 'make' against several tens of thousands of files. The file attributes are checked (metadata operations) and if necessary, the file is read, compiled, then data is written back out to storage.

2.2.4.1.1 SWBUILD File Operation Distribution

	Operation	%	Operation	%
	read	0	read file	6
Ę	mmap read	0	rand read	0
ltio	write	0	write file	7
ribı	mmap write	0	rand write	0
Dist	rmw	0	append	0
n [mkdir	1	rmdir	0
atio	readdir	2	create	1
per	unlink	2	unlink2	0
File O	stat	70	access	6
	rename	0	copyfile	0
	locking	0	chmod	5
	statfs	0	pathconf	0

	Slot	Start	End	%
	0	1	511	1
	1	512	1023	5
_	2	1024	2047	7
tion	3	2048	4095	7
ipni	4	4096	4096	0
istri	5	4096	8191	45
۵ ۵	6	8192	8192	0
Siz	7	8192	16383	13
sfer	8	16384	16384	0
rans	9	16384	32767	3
Р	10	32768	65535	2
Rea	11	65536	65536	0
	12	98304	98304	0
	13	65536	131072	17
	14	262144	262144	0
	15	524288	524288	0

2.2.4.1.2 SWBUILD Read Transfer Size Distribution

2.2.4.1.3 SWBUILD Write Transfer Size Distribution

	Slot	Start	End	%
	0	1	511	5
	1	512	1023	3
_	2	1024	2047	10
tion	3	2048	4095	15
ibu	4	4096	4096	0
istr	5	4096	8191	14
e D	6	8192	8192	0
Siz	7	8192	16383	7
sfei	8	16384	16384	0
Lan	9	16384	32767	6
e T	10	32768	65535	4
Vrit	11	65536	131072	36
>	12	98304	98304	0
	13	131072	131072	0
	14	262144	262144	0
	15	524288	524288	0

	Option	Value	Option	Value
Miscellaneous	write commit %	33	background	0
	% direct	0	sharemode	0
	% osync	0	uniform size dist	0
	% notification	0	init rate throttle	0
	LRU	0	init read flag	1
	release version	1		

2.2.4.1.4 SWBUILD Miscellaneous Characteristics

2.2.4.1.5 SWBUILD Access Pattern Characteristics

rns	Option	Value	Option	Value
ttei	rand dist behavior	0	% per spot	0
s Pa	min acc per spot	0	access mult spot	5
cess	affinity %	0	spot shape	0
Acc	geometric %	10	align	0

2.2.4.1.6 SWBUILD Content Pattern Characteristics

SU	Option	Value	Option	Value
ter	dedup %	0	dedup within %	0
Pat	dedup across %	0	dedup group count	1
ent	dedup granule size	4096	dedup gran rep limit	100
nte	compress %	80	comp granule size	8192
ပိ	cipher flag	0	pattern version	1

2.2.4.1.7 SWBUILD Execution Parameters

Execution Parameters	Parameter	Value	Parameter	Value
	Procs	5	Dirs per proc	50
	Oprate per proc	100	Files per dir	100
	Avg file size			16 KiB

2.2.4.1.8 SWBUILD Validation Criteria

Threshold	Value
Per proc oprate	>= 75% of requested
Global average oprate	>= 95% of requested
Per proc maximum acceptable latency	N/A
Global average maximum acceptable latency	N/A
Workload oprate ratio variance	N/A

2.2.5 Video Data Acquisition (VDA) Benchmark

2.2.5.1 Workload description

The workload generally simulates applications that store data acquired from a temporally volatile source (e.g. surveillance cameras). A stream refers to an instance of the application storing data from a single source (e.g. one video feed). The storage admin is concerned primarily about maintaining a minimum fixed bit rate per stream and secondarily about maintaining the fidelity of the stream. The goal of the storage admin is to provide as many simultaneous streams as possible while meeting the bit rate and fidelity constraints.

The business metric for the benchmark is STREAMS. The benchmark consists of two workload objects: VDA1 (data stream) and VDA2 (companion applications). Each stream corresponds to a roughly 36 Mb/s bit rate, which is in the upper range of high definition video.

2.2.5.2 Workload characteristics

2.2.5.2.1 VDA File Operation Distribution

	Operation	%	Operation	%
Ę	read	0	read file	0
	mmap read	0	rand read	0
utic	write	100	write file	0
ribı	mmap write	0	rand write	0
Dist	rmw	0	append	0
on [mkdir	0	rmdir	0
File Operatic	readdir	0	create	0
	unlink	0	unlink2	0
	stat	0	access	0
	rename	0	copyfile	0
	locking	0	chmod	0
	statfs	0	pathconf	0

2.2.5.2.1.1 VDA1 File Operation Distribution

	Operation	%	Operation	%
n	read	5	read file	0
	mmap read	0	rand read	84
uti	write	0	write file	0
trib	mmap write	0	rand write	0
Dist	rmw	2	append	0
u o	mkdir	0	rmdir	0
rati	readdir	3	create	1
bei	unlink	1	unlink2	0
e O	stat	2	access	2
Ξ	rename	0	copyfile	0
	locking	0	chmod	0
	statfs	0	pathconf	0

2.2.5.2.1.2 VDA2 File Operation Distribution

2.2.5.2.2 VDA Read Transfer Size Distribution

2.2.5.2.2.1 VDA1 Read Transfer Size Distribution

	Slot	Start	End	%
	0	1	511	0
	1	512	1023	0
_	2	1024	2047	0
tio	3	2048	4095	0
ibu	4	4096	4096	0
istri	5	4097	8191	0
ē	6	8192	8192	0
Siz	7	8193	16383	0
sfer	8	16384	16384	0
rans	9	16385	32767	0
Р	10	32768	32768	0
Rea	11	65536	65536	15
	12	131072	131072	10
	13	262144	262144	20
	14	524288	524288	35
	15	1048576	1048576	20

	Slot	Start	End	%
	0	1	511	0
	1	512	1023	0
c	2	1024	2047	0
fi	3	2048	4095	0
ribu	4	4096	4096	0
isti	5	4097	8191	0
e D	6	8192	8192	0
Siz	7	8193	16383	0
sfer	8	16384	16384	0
ans	9	16385	32767	0
Ĕ,	10	32768	32768	0
lead	11	65536	65536	15
•	12	131072	131072	10
	13	262144	262144	20
	14	524288	524288	35
	15	1048576	1048576	20

2.2.5.2.2. VDA2 Read Transfer Size Distribution

2.2.5.2.3 VDA Write Transfer Size Distribution

2.2.5.2.3.1 VDA1 Write Transfer Size Distribution

	Slot	Start	End	%
	0	1	511	0
	1	512	1023	0
e	2	1024	2047	0
tion	3	2048	4095	0
ibu	4	4096	4096	0
istr	5	4097	8191	0
e D	6	8192	8192	0
· Siz	7	8193	16383	0
sfei	8	16384	16384	0
ran	9	16385	32767	0
ЕT	10	32768	32768	5
Vrit	11	65536	65536	10
	12	131072	131072	10
	13	262144	262144	25
	14	524288	524288	25
	15	1048576	1048576	25

	Slot	Start	End	%
	0	1	511	0
	1	512	1023	0
Ę	2	1024	2047	0
ltio	3	2048	4095	0
ribu	4	4096	4096	0
ist	5	4097	8191	0
e 🛛	6	8192	8192	0
Siz	7	8193	16383	0
sfei	8	16384	16384	0
ran	9	16385	32767	0
еŢ	10	32768	32768	5
Vrit	11	65536	65536	10
5	12	131072	131072	10
	13	262144	262144	25
	14	524288	524288	25
	15	1048576	1048576	25

2.2.5.2.4 VDA Miscellaneous Characteristics

2.2.5.2.4.1 VDA1 Miscellaneous Characteristics

	Option	Value	Option	Value
ns	write commit %	5	background	0
Jeo	% direct	0	sharemode	0
scellar	% osync	0	uniform size dist	0
	% notification	0	init rate throttle	0
Ξ	LRU	0	init read flag	1
	release version	1		

2.2.5.2.4.2 VDA2 Miscellaneous Characteristics

	Option	Value	Option	Value
SUC	write commit %	0	background	0
Jec	% direct	0	sharemode	0
la	% osync	0	uniform size dist	0
sce	% notification	0	init rate throttle	0
Σ	LRU	0	init read flag	1
	release version	1		

2.2.5.2.5 VDA Access Pattern Characteristics

2.2.5.2.5.1 VDA1 Access Pattern Characteristics

rns	Option	Value	Option	Value
ittei	rand dist behavior	0	% per spot	0
S Pa	min acc per spot	0	access mult spot	5
ces	affinity %	0	spot shape	0
Ac	geometric %	0	align	0

2.2.5.2.5.2 VDA2 Access Pattern Characteristics

		Option	Value	Option	Value
Access Patterns	ns	rand dist behavior	0	% per spot	0
	ttei	min acc per spot	0	access mult spot	5
	Pai	affinity %	0	spot shape	0
		geometric %	0	align	0

2.2.5.2.6 VDA Content Pattern Characteristics

2.2.5.2.6.1 VDA1 Content Pattern Characteristics

Su	Option	Value	Option	Value
ter	dedup %	0	dedup within %	0
Pat	dedup across %	0	dedup group count	1
ent	dedup granule size	4096	dedup gran rep limit	100
nte	compress %	0	comp granule size	8192
ö	cipher flag	0	pattern version	1

2.2.5.2.6.2 VDA2 Content Pattern Characteristics

ns	Option	Value	Option	Value
ter	dedup %	0	dedup within %	0
Pat	dedup across %	0	dedup group count	1
ent	dedup granule size	4096	dedup gran rep limit	100
onte	compress %	0	comp granule size	8192
ŭ	cipher flag	0	pattern version	1

2.2.5.2.7 VDA Execution Parameters

2.2.5.2.7.1 VDA1 Content Pattern Characteristics

kecution rameters	Parameter	Value Parameter	Value
	Procs	1 Dirs per proc	1
	Oprate per proc	9 Files per dir	1
Pa E	Avg file size		1 GiB

			item i attern Characterist	105
n ers	ers	Parameter	Value Parameter	Val
ltio	lete	Procs	1 Dirs per proc	1

2.2.5.2.7.2 VDA2 Content Pattern Characteristics

2.2.5.2.8 VDA Validation Criteria

Oprate per proc

Avg file size

Threshold	Value
Per proc oprate	>= 75% of requested
Global average oprate	>= 95% of requested
Per proc maximum acceptable latency	N/A
Global average maximum acceptable latency	N/A
Workload oprate ratio variance	<= +/- 5%

1

1 GiB

2.2.6 Virtual Desktop Infrastructure (VDI) Benchmark

1 Files per dir

2.2.6.1 Workload description

This workload simulates a steady-state high-intensity knowledge worker in a VDI environment that uses full clones. This workload does not simulate a linked-clone environment. This is the behavior that was seen in traces between the hypervisor and storage when the VM's were running on ESXi, Hyper-V, KVM and Xen environments.

2.2.6.2 Workload characteristics

2.2.6.2.1 VDI File Operation Distribution

	Operation	%	Operation	%
c	read	6	read file	0
	mmap read	0	rand read	20
utio	write	9	write file	0
ribı	mmap write	0	rand write	64
Dist	rmw	0	append	0
e Operation I	mkdir	0	rmdir	0
	readdir	0	create	0
	unlink	0	unlink2	0
	stat	0	access	1
Ē	rename	0	copyfile	0
	locking	0	chmod	0
	statfs	0	pathconf	0

	Slot	Start	End	%
	0	512	512	1
	1	2048	2048	1
_	2	2560	3584	1
tion	3	4096	4096	20
ibu	4	4608	7680	1
istri	5	8192	8192	4
۵ ۵	6	8704	15872	4
Siz	7	16384	16384	42
sfer	8	16896	32256	3
rans	9	32768	32768	14
Ер	10	33280	64024	1
Rea	11	65536	65536	6
	12	66048	126976	1
	13	131072	131072	1
	14	262144	262144	0
	15	524288	524288	0

2.2.6.2.2 VDI Read Transfer Distribution

2.2.6.2.3 VDI Write Transfer Distribution

	Slot	Start	End	%
	0	512	512	21
	1	1024	1024	2
-	2	2048	2048	1
tion	3	4096	4096	47
ibu	4	8192	8192	6
istr	5	8704	15360	3
e D	6	16384	16384	5
· Siz	7	16896	30720	5
sfer	8	32768	32768	1
Lan	9	35328	64000	3
e T	10	65536	65536	2
Vrit	11	69632	126976	1
>	12	131072	131072	3
	13	262144	262144	0
	14	524288	524288	0
	15	1048576	1048576	0

2.2.6.2.4 VDI Miscellaneous Characteristics

	Option	Value	Option	Value
S N	write commit %	100	background	0
Jeo	% direct	100	sharemode	0
la	% osync	0	uniform size dist	0
sce	% notification	0	init rate throttle	0
Σ	LRU	0	init read flag	1
	release version	1		

2.2.6.2.5 VDI Access Pattern Characteristics

ns	Option	Value	Option	Value
ttei	rand dist behavior	0	% per spot	0
, Pa	min acc per spot	0	access mult spot	5
ces	affinity %	0	spot shape	0
Ac	geometric %	90	align	0

2.2.6.2.6 VDI Content Pattern Characteristics

SU	Option	Value	Option	Value
ter	dedup %	0	dedup within %	0
Pat	dedup across %	0	dedup group count	1
ent	dedup granule size	4096	dedup gran rep limit	100
nte	compress %	60	comp granule size	8192
ပိ	cipher flag	0	pattern version	1

2.2.6.2.7 VDI Execution Parameters

on ers	Parameter	Value	Parameter		Value
utic	Procs	2	Dirs per proc		1
(eci	Oprate per proc	100	Files per dir		1
Pa E	Avg file size			50	0 MiB

2.2.6.2.8 VDI Validation Criteria

Threshold	Value
Per proc oprate	>= 75% of requested
Global average oprate	>= 90% of requested
Per proc maximum acceptable latency	N/A
Global average maximum acceptable latency	N/A
Workload oprate ratio variance	N/A

3 Installing and Configuring the Benchmark Environment

This section provides information on hardware/software configuration requirements for the load generators and the storage solutions. It also includes installation instructions for the benchmark on the load generators for each of the supported operating systems.

3.1 Setting up the Solution Under Test (SUT)

There are several things you must set up on your storage solution before you can successfully execute a benchmark run.

 Configure enough disk space. You may mount your test disks anywhere in your server's file name space that is convenient for you. The maximum ops/sec a storage solution can process is often limited by the number of independent disk drives configured on the server. In the past, a disk drive could generally sustain on the order of 100-200 NFS or SMB ops/sec. This was only a rule of thumb, and this value will change as new technologies become available. However, you will need to ensure you have sufficient disks configured to sustain the load you intend to measure. Space requirements scale with the requested load.

DATABASE= 24 GiB per DATABASEEDA= 11 GiB per JOB_SETSWBUILD= 5 GiB per BUILDVDA= 24 GiB per STREAMVDI= 12 GiB per DESKTOP

- 2. Initialize (if necessary or desired) and mount all file systems. According to the Run and Reporting Rules, it is not necessary to (re-)initialize the solution under test prior to a benchmark run. However, in the full disclosure report for a benchmark run, any configuration steps or actions taken since the last (re-)initialization must be documented for each component of the solution. Therefore, it may be desirable to re-initialize the solution between runs, depending on the tester's objective. See section 5.2 "Solution File System Creation and Configuration" in the SPEC SFS® 2014 SP2 Run and Reporting Rules for more detail.
- 3. Export or share all file systems to all clients. This gives the clients permission to mount/map, read, and write to your test storage. The benchmark program will fail without this permission.
- 4. Verify that all RPC services work. The clients may use port mapping, mount, and NFS services, or Microsoft name services, and file sharing, provided by the server. The benchmark will fail if these services do not work for all clients on all networks. If your client systems have NFS client software installed, one easy way to do this is to attempt mounting one or more of the server's exported file systems on the client. On a Windows client one may try mapping the shares to ensure that the services are correctly configured on the SMB server.
- 5. Ensure your solution is idle. Any other work being performed by your solution is likely to perturb the measured throughput and response time. The only safe way to make a repeatable measurement is to stop all non-benchmark related processing on your solution during the benchmark run.
- 6. Ensure that your test network is idle. Any extra traffic on your network will make it difficult to reproduce your results, and will probably make your solution look slower. The easiest thing to do is to have a separate, isolated network for all components that comprise the solution. Results obtained on production networks may not be reproducible. Furthermore, the benchmark may fail to correctly converge to the requested load rate and behave erratically due to varying ambient load on the network. **Please do not run this benchmark over a corporate LAN. It can present heavy loads and adversely affect others on the same shared network**.

At this point, your solution should be ready for a benchmark measurement. You must now set up a few things on your client systems so they can run the benchmark programs.

3.2 Setting up the Load Generators

Running the SfsManager requires that the Python 2.7 be installed.

On UNIX systems, create "spec" user. SPEC SFS 2014 benchmark runs should be done as a non-root user. The SPEC SFS 2014 binary must be installed on clients.

To install the SPEC SFS 2014 binary:

- On all the clients:
 - 1. Login as "root"
 - 2. Change directory to the top level directory containing the SPEC SFS 2014 benchmark files
 - 3. Enter 'python SfsManager -install-dir="destination_directory""

Alternately, just copy the binary file to all clients using any feasible method.

Verify that all clients, including prime, have the SFS 2014 binary at the same location with the same name and that this matches the EXEC_PATH parameter in the RC file.

Additional client setup validation:

- 1. Configure and verify network connectivity between all clients and server. Clients must be able to send IP packets to each other and to the server. How you configure this is system-specific and is not described in this document. Two easy ways to verify network connectivity are to use a "ping" program or the netperf benchmark (*http://www.netperf.org*).
- 2. Before starting the benchmark, ensure that the prime client can execute commands on the remote clients using ssh with no password challenges. Refer to Appendix B for an example of how to do this. (On Unix-based systems)
- 3. Ensure that the file systems specified in the CLIENT_MOUNTPOINTS string are mounted and accessible on all clients. Where possible, it may help to mount all test file systems to a path under a tmpfs directory. This avoids filling local disks if a mount fails. If using Windows clients with UNC paths in CLIENT_MOUNTPOINTS, verify that the clients can access those UNC paths with the USER and PASSWORD specified in the RC file.
- 4. Clients must have free space for logs generated during the course of the run. In general, logs can vary between 500 KiB to 100 MiB or more. If necessary, a separate log path can be specified in the RC file.

* **IMPORTANT** * – If Windows Firewall is turned on; each program will need to be added to the exceptions list. Either open the Windows Firewall control panel and add the applications manually, or wait for the pop-up to appear after the first execution of each application. Other locally-based firewall applications may require a similar allowance. Note that on Windows systems additional firewall options appear, and may default to enabled, upon joining a domain. After joining a domain, verify that your firewall settings are as expected.

* **IMPORTANT** * – Windows client load generator configurations must have one additional client that is used as the Prime client and this client cannot be used to generate load. This constraint is due to Windows security mechanisms that prevent a client from logging into itself. You may use a single client on non-Windows clients, but it is recommended that the prime client not generate load, which would require at least two clients.

3.2.1 Configuring SPEC SFS 2014 Windows Clients for Auto-Startup

The following are the steps to follow to configure Windows clients in order to allow the Prime Client to communicate with them directly and remotely start the SfsManager process when a benchmark run is started.

Granting DCOM Remote Launch permissions:

- 1. Click Start, click Run, type DCOMCNFG, and then click OK.
- 2. In the Component Services dialog box, expand Component Services, expand Computers.
- 3. Right mouse click on My Computer and select properties. **The My Computer dialog box appears.**
- 4. In the My Computer dialog box, click the COM Security tab.
- 5. Under Launch and Activate Permissions, click Edit Limits.
- 6. In the Launch Permission dialog box, follow these steps if your name or your group does not appear in the Groups or user names list:
 - a. In the Launch Permission dialog box, click Add.
 - b. In the Select Users, Computers, or Groups dialog box, add your name and the group in the Enter the object names to select box, and then click OK.
- 7. In the Launch Permission dialog box, select your user and group in the Group or user names box. In the Allow column under Permissions for User, select Remote Launch, and then click OK.

3.3 Configuring Benchmark Parameters

Once you have the clients and server configured, you must set some parameters for the benchmark itself, which you do in a file called the "sfs_rc file". The actual name of the file is a prefix picked by you, and the suffix "_rc". The default version shipped with the benchmark is delivered as "sfs_rc" in the benchmark source directory. One may use any text editor to modify parameters in the rc files.

There are several parameters you must set, and several others you may change to suit your needs while performing a disclosable run. There are also many other parameters you may change which change the benchmark behavior, but lead to a non-disclosable run (for example, turning on the PIT_SERVER). See the SPEC SFS 2014 Run Rules for the classification of all the parameters.

The parameters you must set are:

- 1. **BENCHMARK:** The name of the workload to run. DATABASE, EDA, SWBUILD, VDA, or VDI.
- 2. CLIENT_MOUNTPOINTS: This parameter specifies the names of the file systems the clients will use when testing the storage solution. The business metric values are spread among the client mount points in the following way. If the number of items N in the CLIENT_MOUNTPOINTS list is greater than the business metric value L (the current value for LOAD), then the first L items from the list are used, one business metric value per client/mountpoint. If L>N, then the N+1 business metric value will wrap around to the beginning of the list and allocation proceeds until all L business metrics have been allocated, wrapping around to the beginning of the list as many times as is necessary.

Examples:

For an NFS configuration: client_name:/mount_point_path client_name:/mount_point_path

For a SMB configuration: (client_name followed by UNC path) client_name:\\server\path client_name:\\server\path ...

When using an external file: (CLIENT_MOUNTPOINTS=mountpoints.txt), the syntax for each line in the file is "client_name path". Multiple mountpoints for a single load generator can be specified on the same line, separated by spaces. However, because the CLIENT_MOUNTPOINTS list is used in order, this may create artificial bottlenecks if not done carefully. The lines do not need to be unique. For example:

client1 /mnt client1 /mnt client2 /mnt client3 /mnt1 client3 /mnt2

- **Reminder**: If using Windows load generators, the Prime client must not be listed in the CLIENT_MOUNTPOINTS list.
 - 3. LOAD, INCR_LOAD, and NUM_RUNS: These parameters specify the aggregate load the clients will generate. To test a set of evenly spaced load points, set all three parameters. Set LOAD to the lowest load level, set INCR_LOAD the amount you would like to increase the load for each measured run, and set NUM_RUNS to the number of times you would like to increment the load. This is the easiest way to configure a disclosable run. For example, if you would like to measure 10 evenly spaced points ending at 2000, you would set LOAD to 200, INCR_LOAD to 200 and NUM_RUNS to 10.

When choosing LOAD values please take into consideration the amount of RAM that these workloads will consume on the client.

DATABASE	= 55 MiB per DATABASE
EDA	= 12 MiB per JOB_SET
SWBUILD	= 400 MiB per BUILD
VDA	= 10 MiB per STREAM
VDI	= 8 MiB per DESKTOP
Also, please con	sider the capacity requirements for each LOAD increment
DATABASE	= 24 GiB per DATABASE
EDA	= 11 GiB per JOB_SET
SWBUILD	= 5 GiB per BUILD
VDA	= 24 GiB per STREAM
VDI	= 12 GiB per DESKTOP

- 4. **EXEC_PATH**: Set this to the absolute path to the benchmark executable. The same path will be used on all clients, so the executable must be at the same path on all clients.
- 5. **USER**: Set this to the User ID for launching the benchmark on all clients. (On Windows systems this includes the Domain/User) E.g. DOMAIN\User33
- 6. **PASSWORD**: Set this to the account password for running the benchmark. (Windows clients only)

3.3.1 Other Variables in the RC File

In addition to the parameters required to be changed for a run, the following parameters are optionally adjustable in the RC file – note that some may not be changed or set for a publishable run:

- PRIME_MON_SCRIPT and PRIME_MON_ARGS: This is the name (and argument list) of a
 program which the SPEC SFS 2014 benchmark will execute during the various phases of the
 benchmark. This is often used to start some performance measurement program while the
 benchmark is running to assist with debugging and tuning your system.
 An example monitor script sfs_ext_mon is provided in the SPEC SFS 2014 source directory.
 For a disclosable run, this program/script must be performance neutral and its actions must comply
 with the SPEC SFS® 2014 SP2 Run and Reporting Rules. If this option is used, the script used, as
 well as the contents of the PRIME_MON_ARGS parameter, must be disclosed in the Other
 Solution Notes (otherSutNotes) field. Longer scripts may be attached as a configuration diagram
 and referenced in the Other Solution Notes (otherSutNotes) field.
- 2. **WARMUP_TIME**: Sets the duration of the WARMUP phase in seconds. This may be set to a value between 300 seconds and one week (604,800 seconds.)
- 3. **RUNTIME***: Sets the duration of the RUN (measurement) phase of the benchmark.
- 4. WORKLOAD_FILE*: Used to import custom workload objects.
- 5. **IPV6_ENABLE:** Flag to set for when the benchmark should use IPv6 to communicate with other benchmark processes.
- 6. **INIT_RATE:** Set dataset creation rate throttle such that no proc will exceed this rate in MiB/s during the INIT phase
- 7. **CLIENT_MEM*:** Used to tell the benchmark the amount of RAM in the client. This is used to set the file size the benchmark will use if the FILE_SIZE parameter is not set in the RC or benchmarks.xml file. This parameter has no effect for any publishable workloads.
- 8. AGGR_CAP*: Used to tell the benchmark the maximum aggregate data set size. This is used to set the file size the benchmark will use if the FILE_SIZE parameter is not set in the RC or benchmarks.xml file. This parameter has no effect for any publishable workloads.
- 9. **FILE_SIZE*:** Used to tell the benchmark to create files of a particular size size. This overrides the effects of CLIENT_MEM and AGGR_CAP, both of which are ignored if this is set. This is overridden by the value set in benchmarks.xml for the requested workload.
- 10. **DIR_COUNT*:** Used to specify the number of subdirectories at the leaf level. This is overridden by the value set in benchmarks.xml for the requested workload.
- 11. **FILES_PER_DIR*:** Used to specify the number of files in each leaf subdirectory. This is overridden by the value set in benchmarks.xml for the requested workload.
- 12. UNLINK_FILES*: Specify whether to delete the dataset between load points. This is not recommended for common test cases.
- 13. LATENCY_GRAPH*: Specify whether to or not to enable timers to keep track of latency per op type. Enabling this prints out a table of op type latencies and overall latency histogram for each iteration of the benchmark. If this value is not present in the RC file, the behavior will be enabled by SfsManager.
- 14. **HEARTBEAT_NOTIFICATIONS*:** Specify whether or not to display heartbeat notifications during all phases of benchmark execution to indicate ongoing test activity. If this value is not present in the RC file, the behavior will be enabled by SfsManager.
- 15. **DISABLE_FSYNCS*:** Disables fsync() calls in the load generator.
- 16. USE_RSHRCP*: Use RSH instead of SSH to start remote processes on UNIX platforms.
- 17. **BYTE_OFFSET*:** Over-ride beginning offset of files adds a fixed-length offset to change the beginning offset of every file.
- 18. MAX_FD*: Sets the maximum number of file descriptors each proc can use.
- 19. **PIT_SERVER*:** Name of the server running the Programmable Interval Timer server.
- 20. **PIT_PORT*:** TCP port number to access the PIT on the remote PIT server.
- 21. LOCAL_ONLY*: Use sh instead of ssh/rsh confines the benchmark to run only on the prime client. This works on UNIX and Windows systems. This option is for benchmark development

```
SPEC SFS® 2014 SP2
```

testing purposes only and not recommended for use. This option may be deprecated in future releases.

- 22. **FILE_ACCESS_LIST*:** If enabled, the benchmark will dump a list of all files accessed during each load point.
- 23. **SHARING_MODE*:** If enabled, the dataset for a single business metric is shared between all procs in that business metric.
- 24. **SOCK_DEBUG*:** Enables TCP socket debug logging.
- 25. TRACEDEBUG*: Enables tracing of file system operations for each proc.
- 26. NO_SHARED_BUCKETS*: Disable creation of dataset for unused op types and sharing of dataset between op types.
- 27. NO_OP_VALIDATE*: Disable validation that all op types can succeed on file systems under test before each load point.
- 28. UNLINK2_NO_RECREATE*: Suppress recreation of files (with non-zero size) after unlink2.
- 29. **NETMIST_LOGS:** Used to set a non-default location for the netmist_C*.log files. If this isn't set either /tmp/ or c:\tmp\ will be used.

*This parameter may not be changed or set to a non-default value for a publishable run.

4 Running the Benchmark and Interpreting Results

This section contains information on the SPEC SFS 2014 benchmark directory structure, running the benchmark, and interpreting the benchmark metrics output generated in the summary results file.

4.1 SFS Benchmark Directory Structure

The following is a quick overview of the benchmark's directory structure. Please note that the variable "\$SPEC" used below represents the full path to the install_directory, where the benchmark is installed.

1. \$SPEC

The directory contains the SPEC SFS 2014 benchmark Makefile. The makefile is used to build tools, compile the benchmark source into executables, and to clean directories of all executables. Pre-built binaries are provided for many operating systems, therefore compilation is probably not required. The top-level directory also contains the SfsManager and SpecReport tools as well as the example sfs_rc and sfs_ext_mon files.

2. \$SPEC/bin

The benchmark binaries for the specific environment being used are located in the "\$SPEC/bin" directory if the user has built the binaries using the Makefile provided.

3. \$SPEC/binaries

Contains the pre-built binaries for various operating systems.

- 4. \$SPEC/docs Contains documentation for the SPEC SFS 2014 benchmark.
- \$SPEC/msbuild Contains the Microsoft Visual Studio solution file used to compile the benchmark on Windows.
- 6. \$SPEC/netmist

Contains the source files for the netmist load generator.

- \$SPEC/redistributable_sources This directory contains tools relevant to the execution or analysis of SPEC SFS 2014 benchmark runs licensed under compatible terms.
- 8. \$SPEC/win32lib
- Contains compatibility libraries for building the benchmark under Microsoft Visual Studio.
- 9. \$SPEC/results

Contains benchmark log and results files created during a benchmark run. This directory is created upon successful start of benchmark execution if it does not exist.

4.2 Pre-Compiled SPEC SFS 2014 Benchmark Binaries

The SPEC SFS 2014 benchmark includes pre-compiled binaries for a large number of supported platforms and architectures. If it is necessary or desired for the user to compile a version of the benchmark source for testing, a generic UNIX makefile is provided in the benchmark top level directory (\$SPEC). The makefile may be modified or supplemented in a performance neutral fashion to facilitate the compilation and execution of the benchmark on operating systems not included within the benchmark distribution. To build the software simply type: make

The Visual Studio solution file is also provided to compile the Windows executables. The solution file is located in the PEC/msbuild subdirectory. The SPEC SFS 2014 benchmark can be built with Visual Studio C++ 2015 Express.

The following is a list of the vendors and their respective operating system levels for which the benchmark workloads have been pre-compiled and included with the benchmark distribution.

IBM Corporation AIX 7.2 FreeBSD FreeBSD 11 Oracle Corporation Solaris 11.1, Solaris 11.3 Red Hat, Inc. RHEL 6, RHEL 7 Apple Inc. Mac OS X Microsoft Corporation Windows 7, Windows 8, Windows 10, Windows Server 2008R2, Windows Server 2012, Windows Server 2012R2, Windows Server 2016

4.3 Using the SFS Manager

The SfsManager is used to run the benchmark. The results obtained from multiple data points within a run are also collected in a form suitable for use with other result formatting tools.

4.3.1 Example of SUT Validation

By default, and during a publication run, the client validates that it can perform all of the POSIX level operations that will be used during the benchmark before starting benchmark execution. If the validation fails, then the benchmark will terminate, with errors collected in the log files.

4.3.2 Example of a Benchmark Run

```
<<< Thu Oct 12 06:57:11 2017: Starting VDI run 1 of 1: DESKTOPS=1 >>
```

```
SPEC SFS2014 SP2 Release
                                $Revision: 1688 $
 This product contains benchmarks acquired from several sources who
 understand and agree with SPEC's goal of creating fair and objective
 benchmarks to measure computer performance.
 This copyright notice is placed here only to protect SPEC in the
 event the source is misused in any manner that is contrary to the
 spirit, the goals and the intent of SPEC.
 The source code is provided to the user or company under the license
 agreement for the SPEC Benchmark Suite for this product.
_____
This program contain a 64-bit version of Mersenne Twister pseudorandom
number generator.
Copyright (C) 2004, Makoto Matsumoto and Takuji Nishimura,
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modification, are permitted provided that the following conditions
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```

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Test run time = 300 seconds, Warmup = 300 seconds. Op latency reporting activated Files per directory set to 1 Directories per proc set to 1 Using custom file size 51200 Kbytes Running 2 copies of the test on 1 clients Results directory: /home/capps/Downloads/trunk/bin/results Clients have a total of 1024 MiBytes of memory Clients have 512 MiBytes of memory size per process Clients each have 2 processes Adjustable aggregate data set value set to 1024 MiBytes Each process file size = 51200 kbytes Client data set size = 1100 MiBytes Total starting data set size = 1100 MiBytes Total initial file space = 1100 MiBytes Total max file space = 1200 MiBytes

[INFO][Thu Oct 12 06:57:11 2017]Exec validation successful

SPEC SFS2014_SP2 Release \$Revision: 1689 \$

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This program contain a 64-bit version of Mersenne Twister pseudorandom number generator.

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 Band 1:
 20us:7
 40us:257
 60us:163
 80us:96

 Band 2:
 200us:75
 400us:319
 600us:2973
 800us:4387

 Band 3:
 2ms:15518
 4ms:26123
 6ms:5918
 8ms:230

 Band 4:
 12ms:60
 14ms:56
 16ms:53
 18ms:30

 4ms:26123
 6ms:5918
 8ms:230

 14ms:56
 16ms:53
 18ms:30
 10ms:103 20ms:17

Band 5:	40ms:19	60ms:0	80ms:0	100ms:0	
Band 6:	200ms:0	400ms:0	600ms:0	800ms:0	1s:0
Band 7:	2s:0	4s:0	6s:0	8s:0	10s:0
Band 8:	20s:0	40s:0	60s:0	80s:0	120s:0
Band 9:	120+s:0				

netmist completed successfully, summarizing.

Reminder: The benchmark run may take many hours to complete depending upon the requested load and how many data points were requested. Also, some failures may take more than an hour to manifest.

5 Submission and Review Process

The SPEC SFS 2014 benchmark release includes a tool for collecting benchmark results in a format that can be submitted by email to the SPEC SFS 2014 results processing facility at SPEC. This facility will automatically process these results and distribute them to the SPEC Storage subcommittee for review. This section describes how you can use these tools to generate a file for each result that you wish to submit to SPEC. It also describes the review process that occurs once the results are submitted. At this point, it is expected that you have become familiar with the SPEC SFS® 2014 SP2 Run and Reporting Rules. See the SPEC SFS® SP2 Run and Reporting Rules documentation that is included in the distribution.

5.1 Creating Reports

Once a benchmark run is completed, the configuration file, results file and additional information are combined into a submission file that is used for submitting runs to SPEC for review using the SpecReport command. Descriptions of the fields that need to be filled out in the submission file are included in Section 6.1 in the SPEC SFS® 2014 SP2 Run and Reporting Rules. This same submission file can be used to generate reports in the form presented on the SPEC web site using the SpecReport command. Each command is documented below.

\$ python SpecReport -h

Usage: python SpecReport [options]

Command Line Option	Description	Required/Optional
[-i <file>] or</file>	Specify XML submission file	Required
[submission-file= <file>]</file>		
[-r <file>] or</file>	Specify RC file	Required for initial
[rc-file= <file>]</file>		package creation
[-s <suffix>] or</suffix>	Suffix used in log and summary files, similar to	Required for initial
[suffix= <suffix>]</suffix>	SfsManager	package creation
[-p <prefix>] or [prefix=<prefix>]</prefix></prefix>	Prefix common to all submission files that get created. Default during initial submission package creation: sfs2014-YYYYmmdd-HHMM. This parameter is required for renaming existing submissions.	Optional
[-u] or [update]	Update an existing submission. This option gets the prefix from the submission file (-i <file>) filename. The RC file, suffix, and results directory flags will be ignored. Use with -p <prefix> for a renamed and updated version of a submission.</prefix></file>	Optional
[-d <dir>] or</dir>	Results directory, default is "results" in the	Optional
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[results-dir= <dir>]</dir>	current working directory.	
[-o <file>] or</file>	Output ZIP archive for full disclosure.	Optional
[output= <air>]</air>		
[-a <file1,file2,>] or</file1,file2,>	List of extra files to attach to the submission (e.g.	Optional
[client/mountpoint file)	
attachments= <file1,file2,>]</file1,file2,>	• · ·	
[validate-only]	Validate the submission without creating the full	Optional
	disclosure package.	
[-h]	Show usage info.	Optional

5.1.1 Creating the Submission Package

To create a submission file one must first create an XML document based on the submission_template.xml example found in the base directory. The template document has the correct XML structure expected by SpecReport. Valid entries for each field can be found in the SPEC SFS® 2014 SP2 Run and Reporting Rules. Edit a copy of the template and fill in each field according to the run rules with specific information pertaining to the SUT.

Once the XML submission document is complete a formal submission package can be created with SpecReport. The tool has 3 required arguments: the RC file, the XML submission file, and the suffix used during the test, the same suffix used with SfsManager. To test the submission files for correctness, issue the command

\$ python SpecReport -r <RC file> -i <XML file> -s <suffix> --validate-only

The tool will check for the existence of all the necessary files and check the format of the XML document. If the command returns without reporting any errors, repeat the command without the "--validate-only" flag to create the submission package. The package will be a zip archive containing the following files: The RC file, the run summary files, the submission file, an HTML version of the report, a text version of the report, a SUB version of the report, and any configuration diagrams.

The syntax for updating an existing set of submission files is

\$ python SpecReport -u -i <XML file>

5.2 Submitting Results

Once you have generated a submission file as described in the Creating the Submission Package section above, you may submit your run for review by the SPEC Storage subcommittee by emailing the ZIP file to subsfs2014@spec.org. Upon receipt, the SPEC results processing facility will parse the submission file and validate the formats. If the check passes, an email reply is returned to the sender including a submission number assigned to the result. This submission number is used to track the result during the review and publishing process. If there are any formatting errors, the parser will respond with a failure message indicating where in the file the parsing failed. You may then either correct the error and resubmit or contact the SPEC office for further assistance.

Every results submission goes through a minimum two-week review process, starting on a scheduled SPEC Storage sub-committee conference call. During the review, members of the committee may contact the submitter and request additional information or clarification of the submission. Once the result has been reviewed and accepted by the committee, it is displayed on the SPEC web site at http://www.spec.org/.

6 FAQ

6.1 SPEC SFS 2014 Benchmark Press Release

Question 1: Answer:	What is the SPEC SFS 2014 benchmark and how does it compare to other storage solution benchmarks? The SPEC SFS 2014 benchmark is the latest version of the Standard Performance Evaluation Corp.'s benchmark that measures a storage solution throughput and response time. It differs from other file server benchmarks in that it provides a standardized method for comparing performance across different vendor platforms. The benchmark was written to be solution independent and vendor-neutral. Results are validated through peer review before publication on SPEC's public website http://www.spec.org/SPECsfs2014/ >		
Question 2: Answer:	What improvements have been made to the SPEC SFS 2014 benchmark? See the "What's New in SPEC SFS 2014?" document included with the benchmark or on the SPEC SFS 2014 website: <u>http://www.spec.org/sfs2014/</u>		
Question 3: Answer:	How were the SPEC SFS 2014 workloads determined? The SPEC SFS 2014 workloads are based on network traces collected from real environments and input from domain experts in and published documentation of the real- world implementation of the application types being simulated.		
Question 4: Answer:	What are the metrics for the SPEC SFS 2014 benchmark?The SPEC SFS 2014 benchmark has multiple performance measurement metrics:• DATABASE= DATABASES• EDA= JOB_SETS• SWBUILD= BUILDS• VDA= STREAMS• VDI= DESKTOPS		
Question 5: Answer:	What is the correlation between the SPEC SFS 2014 benchmark and the TPC (Transaction Processing Council) and SPC (Storage Performance Council) benchmarks? There is no correlation; the benchmarks present very different workloads on the solutions under test and measure different aspects of solution performance.		
Question 6: Answer:	Is the SPEC SFS 2014 benchmark a CPU-intensive or I/O-intensive benchmark? The SPEC SFS 2014 benchmark is an application-level benchmark that heavily exercises CPU, mass storage and network components. The greatest emphasis is on I/O, especially as it relates to operating and file system software. To obtain the best performance for a system running the SPEC SFS 2014 benchmark, the vendor will typically add additional hardware such as memory, disk controllers, disks, network controllers and buffer cache – as needed in order to help alleviate I/O bottlenecks and to ensure that server CPUs are used fully.		
Question 7: Answer:	For what computing environment is the SPEC SFS 2014 benchmark designed? The benchmark was developed for load-generating clients running UNIX or Windows. The SPEC SFS 2014 benchmark can be used to evaluate the performance of any storage solution, regardless of the underlying environment.		

Question 8:	Can users measure performance for workloads other than the ones provided within the SPEC SES 2014 benchmark?
Answer:	Yes, users can measure their own workloads by making changes to the SPEC SFS 2014 benchmark workload objects file. The SPEC SFS® 2014 User's Guide details how this can be done. Workloads created by users cannot, however, be compared with SPEC SFS 2014 results, nor can they be published in any form, as specified within the SPEC SFS 2014 license.
Question 9: Answer:	To what extent is the server's measured performance within the SPEC SFS 2014 benchmark affected by the client's performance? SPEC has written the SPEC SFS 2014 benchmark to include the effect of client performance on SPEC SFS 2014 results. This is a storage solution benchmark, not a component level benchmark. The aggregate data set sweeps a range that covers the in cache and out of cache cases for the solution. This provides coverage for the real world situations.
Question 10: Answer:	How does SPEC validate numbers that it publishes? Results published on the SPEC Web site have been reviewed by SPEC members for compliance with the SPEC SFS® 2014 SP2 Run and Reporting Rules, but there is no monitoring beyond that compliance check. The vendors that performed the tests and submitted the performance numbers have sole responsibility for the results. SPEC is not responsible for any measurement or publication errors.
Question 11: Answer:	Are the reported SPEC SFS 2014 configurations typical of systems sold by vendors? Yes and no. They are similar to large server configurations, but the workload is heavier than that found on smaller server configurations. SPEC has learned from experience that today's heavy workload is tomorrow's light workload. For some vendors, the configurations are typical of what they see in real customer environments, particularly those incorporating high-end servers. For other vendors, SPEC SFS 2014 configurations might not be typical.
Question 12: server?	Do the SPEC SFS® 2014 SP2 Run and Reporting Rules allow results for a clustered
Answer:	Yes, cluster configurations are allowed as long as they conform to the SPEC SFS® 2014 SP2 Run and Reporting Rules.
Question 13: Answer:	What resources are needed to run the SPEC SFS 2014 benchmark? In addition to a server, a test bed includes several clients and an appropriate number of networks. Ideally, the server should have enough memory, disks and network hardware to saturate the CPU. The test bed requires at least one network. Examples of typical load- generating configurations can be found on the SPEC Web site: < http://www.spec.org/SPECsfs2014/>.
Question 14: Answer:	What is the estimated time needed to set up and run the SPEC SFS 2014 benchmark? Hardware setup and software installation time depend on the size of the server and the complexity of the test beds. Many servers require large and complex test beds. The SPEC SFS 2014 software installs relatively quickly. A SPEC SFS 2014 submission from a vendor includes at least 10 data points, with each data point taking from ~30 to ~90 minutes to complete. The performance of the storage solution is a factor in the time it takes to setup and run each load point.

Question 15: Answer:	What shared resources does the SPEC SFS 2014 benchmark use that might limit performance? Shared resources that might limit performance include CPU, memory, disk controllers, disks, network controllers, network concentrators, network switches, clients, etc.
Question 16: Answer:	Does the SPEC SFS 2014 benchmark permit tuning parameters? When submitting results for SPEC review, vendors are required to supply a description of all tuning parameters for all cases where non-default values were used for all components in the SUT. This information is displayed in the appropriate sections in published SPEC SFS 2014 results.
Question 17: Answer:	Can a RAM disk be used within a SPEC SFS 2014 configuration? SPEC enforces strict storage rules for stability. Generally, RAM disks do not meet these rules, since they often cannot survive cascading failure-recovery requirements unless an uninterruptible power supply (UPS) with long survival times is used.
Question 18: Answer:	How will the choice of networks affect SPEC SFS 2014 results? Different link types and even different implementations of the same link type might affect the measured performance for better or worse of a particular server. Consequently, the results measured by clients in these situations might vary as well.
Question 19: Answer:	Is the SPEC SFS 2014 benchmark scalable with respect to CPU, cache, memory, disks, controllers and faster transport media? Yes the benchmark is scalable as users migrate to faster technologies.
Question 20: Answer:	What is the price of a SPEC SFS 2014 license and when will it be available? The SPEC SFS 2014 benchmark is available now from the SPEC download site for US\$2,000. A discounted price is available for non-profit and academic licensees. Contact the SPEC office: (See www.spec.org for any updates) Standard Performance Evaluation Corporation (SPEC) 7001 Heritage Village Plaza Suite 225 Gainesville, VA 20155 Phone: 1-703-579-8460 Fax: 1-703-579-8463 E-Mail: info@spec.org
Question 21: Answer:	Can users get help in understanding how to run the SPEC SFS 2014 benchmark? The majority of questions should be answered in the SPEC SFS® 2014 User's Guide. There is also useful information on the SPEC Web site: <u>http://www.spec.org/SPECsfs2014/</u>
Question 22: Answer:	Do I need to measure every workload? No. Each workload has a separate metric that can be published independently.
Question 23: Answer:	How do I get started running the SPEC SFS 2014 benchmark? Please read the SPEC SFS® 2014 User's Guide and SPEC SFS® 2014 SP2 Run and Reporting Rules in their entirety.
Question 24: Answer:	I am running into problems setting up and running the benchmark. What can I do? The most common problem is usually that file server file systems are not being correctly mounted on the clients. Most of the problems relating to the SPEC SFS 2014 benchmark can be resolved by referring to appropriate sections of the User's Guide, including this FAQ. Other common problems include: hosts file/DNS configuration, firewalls not being
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disabled, password-less SSH not configured, or attempting to run outside a domain in a Windows environment.

Question 25: I have read the SPEC SFS® 2014 User's Guide. But I am still running into problems. What can I do next? Answer: Looking at the sfslog.* and the sfsc* files can give you an idea as to what may have gone wrong. Inspecting the client logs in /tmp/, c:\tmp\, or at NETMIST_LOGS on each load generator can also provide more details on errors. And, as a last resort, you can contact SPEC at support@spec.org. It is assumed that such calls/emails are from people who have read the SPEC SFS® 2014 User's Guide completely, and have met all the prerequisites for setting up and running the benchmark. **Ouestion 26:** How does one abort a run? Answer: The benchmark can be aborted by simply stopping the SfsManager. This will terminate all SFS related processes on all clients and on the prime client. **Ouestion 27:** For a valid run, which parameters are required to be unchanged? Answer: Information is provided in the SPEC SFS® 2014 SP2 Run and Reporting Rules and in the sfs_rc file, and this is enforced by the benchmark. If invalid parameter values are selected, the benchmark reports an invalid run. **Question 28:** Is there a quick way to debug a testbed? Answer: Read the SPEC SFS® 2014 User's Guide, ping the server from the client, ping from the prime client to the other clients and vice versa, validate that the paths in CLIENT MOUNTPOINTS exist on all load generators and are writeable, validate password-less SSH works from the prime client to all load generators, run the benchmark with one client and one file system.

6.2 Tuning the Solution

Question 29:	What are a reasonable set of parameters for running the benchmark?
Answer:	Study existing results pages with configuration information similar to your system configuration.
Question 30:	How do I increase the performance of our solution?
Answer:	One may need to add, as necessary, one or more of the following: processors, memory disks, controllers, etc.

6.3 Submission of Results

Question 31:	We have a valid set of results. How do we submit these results to SPEC?
Answer:	See the Submission and Review Process section above. The new submission tool
	documentation is in that section.

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8 Research corner

8.1 Custom changes to SfsManager to add new workloads.

Each workload has an XML description in the benchmarks.xml file. When adding or changing workloads within the SfsManager, edit this file to define or change the definition of a workload.

```
Example :
```

```
<benchmark name="NEW" business_metric="NEWMETRIC">
    <workload name="WORKLOAD1">
        <oprate>100</oprate>
        <instances>1</instances>
    </workload>
    <workload name="WORKLOAD2">
        <oprate>10</oprate>
        <instances>4</instances>
        <shared files/>
    </workload>
    <dedicated subdirectory/>
    <override_parm name="RUNTIME">300</override_parm>
    <override_parm name="FILE_SIZE">500m</override_parm>
    <override_parm name="DIR_COUNT">10</override_parm>
    <override_parm name="FILES_PER_DIR">100</override_parm>
    <override_parm name="SHARING_MODE">1</override_parm>
    <threshold type="proc oprate">75</threshold>
    <threshold type="global oprate">90</threshold>
    <threshold type="workload variance">10</threshold>
</benchmark>
```

8.2 Custom workload objects

The SPEC SFS 2014 benchmark is capable of running user defined workloads. These are not publishable via SPEC, however consumers may find it useful to create a model of their specific application load and then to be able to run the benchmark while presenting their specific application's workload.

8.2.1 Export current workload definitions

The existing SPEC SFS 2014 workload definitions may be exported to a flat text file by executing the command:

netmist -E > workload.txt

8.2.2 Workload object attribute definitions

Each workload object contains the following attributes:

Workload name	< Name of the Workload object >
Percent Read	< Percent of the Op mix that is sequential read transfers from files >
Percent Read_file	< Percent of the Op mix that is sequential read whoe files >
Percent Mmap_read	< Percent of the Op mix that is using Mmap to read files >
Percent Random_read	< Percent of the Op mix that is reading at random locations >
Percent Write	< Percent of the Op mix that is sequential write transfers to files >
Percent Write_file	< Percent of the Op mix that is sequential write of whole files >
Percent Mmap_write	< Percent of the Op mix that is using Mmap to write files >
Percent Random_write	< Percent of the Op mix that is writing at random locations >
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Percent Read_modify_write < Percent of the Op mix that is read/modify/write at location in a file > Percent Mkdir < Percent of the Op mix that is creating new directories > Percent Rmdir < Percent of the Op mix that is removing existing directories > Percent Unlink < Percent of the Op mix that is removing empty files > < percent of the Op mix that is removing files that have size > Percent Unlink2 < Percent of the Op mix that is creating new files > Percent Create < Percent of the Op mix that is appending writes to existing files > Percent Append Percent Lock < Percent of the Op mix that is locking files > Percent Access < Percent of the Op mix that is calling access() on files >< Percent of the Op mix that is calling stat() on files > Percent Stat < Percent of the Op mix that is calling chmod() on files > Percent Chmod < Percent of the Op mix that is calling readdir() on directories > Percent Readdir < Percent of the Op mix that is copying whole files > Percent Copyfile Percent Rename < Percent of the Op mix that is renaming files > Percent Statfs < Percent of the Op mix that is calling statfs() >< Percent of the Op mix that is calling pathconf() >Percent Pathconf Percent Custom1 < Reserved for future use > Percent Custom2 < Reserved for future use > Read elem 0 xfer min size < Minimum transfer size for this read slot > Read elem 0 xfer max size < Maximum transfer size for this read slot > Read elem 0 xfer percent < Percent of reads that are of this slot definition > Read elem 1 xfer min size < Minimum transfer size for this read slot > Read elem 1 xfer max size < Maximum transfer size for this read slot > Read elem 1 xfer percent < Percent of reads that are of this slot definition > Read elem 2 xfer min size < Minimum transfer size for this read slot > Read elem 2 xfer max size < Maximum transfer size for this read slot > Read elem 2 xfer percent < Percent of reads that are of this slot definition > Read elem 3 xfer min size < Minimum transfer size for this read slot > Read elem 3 xfer max size < Maximum transfer size for this read slot > Read elem 3 xfer percent < Percent of reads that are of this slot definition > Read elem 4 xfer min size < Minimum transfer size for this read slot > Read elem 4 xfer max size < Maximum transfer size for this read slot > Read elem 4 xfer percent < Percent of reads that are of this slot definition > Read elem 5 xfer min size < Minimum transfer size for this read slot > Read elem 5 xfer max size < Maximum transfer size for this read slot > Read elem 5 xfer percent < Percent of reads that are of this slot definition> Read elem 6 xfer min size < Minimum transfer size for this read slot > < Maximum transfer size for this read slot > Read elem 6 xfer max size Read elem 6 xfer percent < Percent of reads that are of this slot definition> Read elem 7 xfer min size < Minimum transfer size for this read slot > Read elem 7 xfer max size < Maximum transfer size for this read slot > Read elem 7 xfer percent < Percent of reads that are of this slot definition> Read elem 8 xfer min size < Minimum transfer size for this read slot > Read elem 8 xfer max size < Maximum transfer size for this read slot > Read elem 8 xfer percent < Percent of reads that are of this slot definition> Read elem 9 xfer min size < Minimum transfer size for this read slot > Read elem 9 xfer max size < Maximum transfer size for this read slot > Read elem 9 xfer percent < Percent of reads that are of this slot definition> Read elem 10 xfer min size < Minimum transfer size for this read slot > Read elem 10 xfer max size < Maximum transfer size for this read slot > Read elem 10 xfer percent < Percent of reads that are of this slot definition> Read elem 11 xfer min size < Minimum transfer size for this read slot > Read elem 11 xfer max size < Maximum transfer size for this read slot > Read elem 11 xfer percent < Percent of reads that are of this slot definition>

Read elem 12 xfer min size Read elem 12 xfer max size Read elem 12 xfer percent Read elem 13 xfer min size Read elem 13 xfer max size Read elem 13 xfer percent Read elem 14 xfer min size Read elem 14 xfer max size Read elem 14 xfer percent Read elem 15 xfer min size Read elem 15 xfer max size Read elem 15 xfer percent Write elem 0 xfer min size Write elem 0 xfer max size Write elem 0 xfer percent Write elem 1 xfer min size Write elem 1 xfer max size Write elem 1 xfer percent Write elem 2 xfer min size Write elem 2 xfer max size Write elem 2 xfer percent Write elem 3 xfer min size Write elem 3 xfer max size Write elem 3 xfer percent Write elem 4 xfer min size Write elem 4 xfer max size Write elem 4 xfer percent Write elem 5 xfer min size Write elem 5 xfer max size Write elem 5 xfer percent Write elem 6 xfer min size Write elem 6 xfer max size Write elem 6 xfer percent Write elem 7 xfer min size Write elem 7 xfer max size Write elem 7 xfer percent Write elem 8 xfer min size Write elem 8 xfer max size Write elem 8 xfer percent Write elem 9 xfer min size Write elem 9 xfer max size Write elem 9 xfer percent Write elem 10 xfer min size Write elem 10 xfer max size Write elem 10 xfer percent Write elem 11 xfer min size Write elem 11 xfer max size Write elem 11 xfer percent Write elem 12 xfer min size Write elem 12 xfer max size Write elem 12 xfer percent Write elem 13 xfer min size Write elem 13 xfer max size Write elem 13 xfer percent

< Minimum transfer size for this read slot > < Maximum transfer size for this read slot > < Percent of reads that are of this slot definition > < Minimum transfer size for this read slot > < Maximum transfer size for this read slot > < Percent of reads that are of this slot definition >< Minimum transfer size for this read slot > < Maximum transfer size for this read slot > < Percent of reads that are of this slot definition > < Minimum transfer size for this read slot > < Maximum transfer size for this read slot > < Percent of reads that are of this slot definition > < Minimum transfer size for this write slot > < Maximum transfer size for this write slot > < Percent of writes that are of this slot definition > < Minimum transfer size for this write slot > < Maximum transfer size for this write slot > < Percent of writes that are of this slot definition > < Minimum transfer size for this write slot > < Maximum transfer size for this write slot > < Percent of writes that are of this slot definition > < Minimum transfer size for this write slot > < Maximum transfer size for this write slot > < Percent of writes that are of this slot definition > < Minimum transfer size for this write slot > < Maximum transfer size for this write slot > < Percent of writes that are of this slot definition > < Minimum transfer size for this write slot > < Maximum transfer size for this write slot > < Percent of writes that are of this slot definition > < Minimum transfer size for this write slot > < Maximum transfer size for this write slot > < Percent of writes that are of this slot definition > < Minimum transfer size for this write slot > < Maximum transfer size for this write slot > < Percent of writes that are of this slot definition > < Minimum transfer size for this write slot > < Maximum transfer size for this write slot > < Percent of writes that are of this slot definition > < Minimum transfer size for this write slot > < Maximum transfer size for this write slot > < Percent of writes that are of this slot definition > < Minimum transfer size for this write slot > < Maximum transfer size for this write slot > < Percent of writes that are of this slot definition > < Minimum transfer size for this write slot > < Maximum transfer size for this write slot > < Percent of writes that are of this slot definition > < Minimum transfer size for this write slot > < Maximum transfer size for this write slot > < Percent of writes that are of this slot definition > < Minimum transfer size for this write slot > < Maximum transfer size for this write slot > < Percent of writes that are of this slot definition >

Write elem 14 xfer min size	< Minimum transfer size for this write slot >
Write elem 14 xfer max size	< Maximum transfer size for this write slot >
Write elem 14 xfer percent	< Percent of writes that are of this slot definition >
Write elem 15 xfer min size	< Minimum transfer size for this write slot >
Write elem 15 xfer max size	< Maximum transfer size for this write slot >
Write elem 15 xfer percent	< Percent of writes that are of this slot definition >
Percent write commit	< Percent of writes that will be flushed/committed >
Percent direct	< Percent of I/O reads and writes that use O_DIRECT >
Align	< Align file accesses to this size boundary for reads and writes >
Percent osync	< Percent of I/O operations that use O_SYNC >
Percent geometric	< Percent of geometric distribution of file accesses >
Percent compress	< Percent of compressibility of file data >
Percent dedup	< Percent of dedupability of the file data >
Percent dedup_within	< Percent of dedupability within any given file >
Percent dedup_across	< Percent of dedupability across files >
Dedupe group count	< Number of dedupe groups that can share across files >
Percent per_spot	< Percent of file size that is used to set the hot spot size >
Min acc_per_spot	< Minimum accesses per hot spot before re-selecting >
Acc mult_spot	< Scaled hits within a hot spot >
Percent affinity	< Affinity to move away from a hot spot >
Spot shape	< Shape of hot spot. 0 = Uniform Rand, 1 = Geometric rand dist >
Dedup Granule size	< Granule size for Dedup. Minimum dedupable unit size >
Dedup gran rep limit	< Maximum number of dedup replicants >
Comp Granule size	< Granule size for compression. Minimum compressible unit size >
Background	< When set does not include in measurement >
Sharemode	< When set, all files and directories are shared across procs >
Uniform size dist	< When set, all files are the same size, else Gaussian dist >
Rand dist behavior	< Distribution pattern for workload within files: 0 = uniform random; 1 = geometric distribution; 2 = geometric across hot spots >
Cipher behavior	< Encrypt all files to create flat character frequency distribution data >
Notification percent	< Percent of writes that will also generate notification events >
LRU	< Enable use of LRU cache for file descriptors >
Pattern version	< Set data fill pattern characteristics: $1 = SP1$, $2 = SP2 >$
Init rate throttle	< Limit init rate per proc to this value in MiB/s >
Init read flag	< Enable read of existing dataset during INIT; 1 = enable >
Release version	< Identify SFS 2014 SP workload version; 1 = SP1; 2 = SP2 >
FS type POSIX	< Default is POSIX. Used for non-POSIX implementations >

8.2.3 Edit workload definitions

The existing workload definitions must be exported to a file to modify the workload definitions to ensure the file format is correct. The contents are in flat ASCII and can be edited with your favorite editor. Do not remove or add any lines to this file. You may only edit the existing lines, and save the file in the same format.

There are currently nine defined workloads: DB_LOG, DB_TABLE, EDA_FRONTEND, EDA_BACKEND, VDA1, VDA2, VDI, SWBUILD, and USERDEF.

The values of each parameter may be modified to create a new workload. Once the modifications are complete, this new workload definition may be imported and used to run the user defined workload. NOTE: The values for the existing workload descriptions may be modified, but the number of workload objects may not be changed.

8.2.4 Import workload definitions

To import a user defined workload, edit the sfs_rc file and modify the value for the WORKLOAD_FILE= parameter to contain the name of the workload definition file that was modified. Example:

WORKLOAD_FILE=workload_file

The custom workload file need only be present on the prime client, just like the rc file.

NOTE: once you define a custom workload, you must define a workload object in the benchmarks.xml file to use the workload with the SfsManager framework. See section 8.1 above for more detail.

8.2.5 Client mountpoint file format

When specifying a separate file for the CLIENT_MOUNTPOINT option in the SPEC SFS 2014 rc file, use the following format (note, additional options are shown for example purposes, but a real mountpoints file would likely be more consistent):

client1 /mnt1 client2 /mnt1 client3 /mnt1 /mnt2 filesize=8192 client4 /mnt1 dircount=5 filesperdir=10 client6 /mnt3 logpath=/mnt1

The file format consists of one load generator hostname or IP followed by one or more mountpoints with optional extra options that are passed to netmist. The file is space-delimited. The available extra options are:

filesize	set size of files in the dataset (size in KiB)
dircount	set the number of directories in the dataset
filesperdir	set the number of files per directory in the dataset
logpath	set a separate log path for this proc

Note: all these options apply for any business metrics (group of procs) assigned to the particular entry in the client mountpoints list. Therefore, in the example above, any business metric assigned to the last entry (client6 /mnt3) would log to the /mnt1 directory on client6.

9 Appendix A – Building SFS Benchmark Components

9.1 Building the SPEC SFS 2014 benchmark for UNIX

Note that it is a requirement that you have a GNU-compatible build environment (e.g., a 'gmake' compatible 'make') in order to build SFS on a UNIX system.

To build SFS, you need to simply:

- *cd* to the top-level spec directory (\$SPEC)
- Type *make*

9.2 Building the SPEC SFS 2014 benchmark for Windows

The SPEC SFS 2014 benchmark for WIN32 was developed using Microsoft Visual Studio 2015 Community. Other versions of the development environment may work for building the benchmark, but these have not been tested.

9.2.1 Update Visual Studio Libraries/Includes

If not already installed, download and install Microsoft Visual Studio 2015 and the Windows SDK. The installation and configuration of Visual Studio for proper operation is beyond the scope of this document.

9.2.2 Open the Visual Studio solution file

Obtain the SPEC SFS 2014 source base and place it in a convenient location on the system to be used for building. In \$SPEC/msbuild there is a Visual Studio solution file that can be used to build the benchmark executables for x86 (32 bit) as well as x64 (64 bit) based systems.

9.2.3 Build the individual project files

If you are familiar with building projects in Visual Studio, you may build the projects as you are accustomed to doing and skip this section.

To build all the projects at once, from the *Build* menu select *Batch Build*. With all projects and configurations selected, click *Build* or *Rebuild All* and the build process will begin. The binaries will be built in *bin\dist\windows\Win32* and/or *bin\dist\windows\x64*.

To build the individual projects, from the *Project* menu select a project under the *Select Active Project* submenu. The selected project will appear in bold in the workspace window. From the *Build* menu, select *Rebuild All* to force a rebuild of all modules in that project, or select *Build [project name]* to build only those modules that have changed or are missing.

Paths and directories within the project settings are all relative to the project file locations, so building the benchmark files should be possible regardless of where the source base is placed on the system.

10 Appendix B – Setting up password-less SSH

Here is a sample script that can be used to set up password-less SSH on Linux clients.

```
# Define the hosts to be involved in the trust here
# DO NOT include the host you are running, it is added by default
hosts="s2 s3 s4 s5 s6"
echo ""
echo ""
echo "This script will generate SSH keys for the specified machines,"
echo " and set up password-less authentication between them."
echo " You will be prompted for passwords several times during this
process."
echo ""
# Get current user
user=`who -m | awk {'print $1'}`
echo "Trust will be configured for user $user"
echo ""
echo "If this is not correct, stop and login as the appropriate user"
echo -n "(RETURN to continue, CTRL-C to exit) "
read continue
# Configure keys on current host
cd $HOME
ssh-keygen -t rsa
cat .ssh/id rsa.pub >> .ssh/authorized keys
chmod 700 .ssh
chmod 600 .ssh/*
for host in $hosts
do
   ssh $user@$host 'ssh-keygen -t rsa'
   ssh $user@$host 'cat .ssh/id rsa.pub' | cat - >>
~/.ssh/authorized keys
done
for host in $hosts
do
  scp .ssh/authorized keys $host:.ssh
   ssh $user@$host 'chmod 700 .ssh ; chmod 600 .ssh/*'
done
exit
```

11 Appendix C – Tunes for the load generators

11.1 Linux tunes

Some or all of the following tuning parameters may be helpful in tuning Linux load generators for optimal performance. The net.core.somaxconn tuning, if available, is highly recommended for any Linux environment running SPEC SFS 2014. Tuning this will avoid excessive delays during benchmark startup/shutdown between iterations.

To make these tunings persistent, use /etc/sysctl.conf – see the documentation for sysctl.conf for more information.

```
if your NIC type == 10 GigE
echo "300000" > /proc/sys/net/core/netdev_max_backlog
fi
echo "131071" > /proc/sys/net/core/rmem_default
echo "131071" > /proc/sys/net/core/rmem_max
echo "131071" > /proc/sys/net/core/wmem_default
echo "131071" > /proc/sys/net/core/wmem_max
echo "4096 87380 8388608" > /proc/sys/net/ipv4/tcp_rmem
echo "4096 87380 8388608" > /proc/sys/net/ipv4/tcp_wmem
echo "128" > /proc/sys/net/core/somaxconn
echo "5536" > /proc/sys/net/core/somaxconn
echo "5" > /proc/sys/net/ipv4/tcp_fin_timeout
```

Power management, and it's impacts on performance: (From an Operating system perspective)

https://docs-old.fedoraproject.org/en-US/Fedora/20/html/Power_Management_Guide/index.html

Limiting RAM in the client

There may be situations where one wishes to limit the amount of RAM that is used in the client. This can be accomplished via kernel boot options either at the boot prompt, or by editing /boot/grub/grub.conf. See the mem=xxxM option in the Linux vendor's installation guide.

11.2 Windows tunes

- Disable power management, or set every possible value to "High performance" mode. <u>https://docs.microsoft.com/en-us/windows-server/administration/performance-tuning/hardware/power/power-performance-tuning</u>
- Disable the Windows firewall on all clients and servers.
 - Note that upon joining a domain, a new firewall setting may appear for domain networks that defaults to enabled
- Disable the Windows Defender, and any other antivirus software, on the clients and servers.
- Disable all slide-show wallpapers, and screen savers.
- Disable automatic updates.
- Disable any interrupt throttling.

Limiting RAM in the client.

• There may be situations where one wishes to limit the amount of RAM that is used in the client. See: <u>http://msdn.microsoft.com/en-us/library/ff542202(v=vs.85).aspx</u> The section on "remove memory" provides information on how to reduce memory.

12 Appendix D – Workload Definition Tables

DB_TA	BLE																	
	Operation	% Operation	%		Option	Value	Option	Value		Slot	Start	End S	%		Slot S	tart I	ind 9	%
	read	1 read file	0		write commit %	100	background	0		0	1	511	0		0	512	512	0
Ę	mmap read	0 rand read	79		% direct	100	sharemode	1		1	512	1023	0		1	1024	1024	0
utio	write	0 write file	0		% osync	0	uniform size dist	1		2	1024	2047	0	_	2	1536	1536	0
rib	mmap write	0 rand write	20	sce	% notification	0	init rate throttle	0	Ĕ	3	2048	4095	0	tio	3	2048	2048	0
Dist	rmw	0 append	0	Σ	LRU	0	init read flag	1	bu	4	4096	4096	0	ibu	4	2560	2560	0
L L	mkdir	0 rmdir	0		release version	1			istr	5	4097	8191	0	istr	5	3072	3072	0
atio	readdir	0 create	0	rns	Option	Value	Option	Value	e	6	8192	8192	99	ēD	6	3584	3584	0
bei	unlink	0 unlink2	0	itte	rand dist behavior	2	% per spot	5	Siz	7	8193	16383	0	Siz	7	4096	4096	0
<u>e</u>	stat	0 access	0	e B	min acc per spot	1000	access mult spot	5	sfer	8	16384	16384	0	sfei	8	4608	4608	0
Ξ	rename	0 copyfile	0	ces	affinity %	20	spot shape	0	an.	9	16385	32767	0	ran	9	5120	5120	0
	locking	0 chmod	0	Ac	geometric %	2	align	0	F.	10	32768	32768	0	e T	10	8192	8192	100
	statfs	0 pathconf	0	ns	Option	Value	Option	Value	Rea	11	65536	65536	0	Vrit	11	12288	12288	0
ds	Threshold	% Threshold	Value	ter	dedup %	0	dedup within %	0		12	98304	98304	0		12	16384	16384	0
hol	proc oprate	75 proc latency	n/a	Pat	dedup across %	0	dedup group count	1		13	131072	131072	0		13	20480	20480	0
res	global oprate	95 global latongy	nla	ent	dedup granule size	4096	dedup gran rep limit	100		14	262144	262144	0		14	24576	24576	0
臣	workload variance	5	11/d	bit	compress %	50	comp granule size	8192		15	1048576	1048576	1		15	32768	32768	0
on ers	Parameter	Value Parameter	Value	ŭ	cipher flag	0	pattern version	1										
utid het	Procs	10 Dirs per proc	1															
xec	Oprate per proc	16 Files per dir	5															
Pa E	Avg file size		200 MiB															

DB_LO	G																	
	Operation	% Operation	%		Option	Value	Option	Value		Slot	Start	End	%		Slot S	tart E	nd १	6
	read	0 read file	0	SUC	write commit %	100	background	0		0	1	511	0		0	512	512	5
5	mmap read	0 rand read	0	neo	% direct	100	sharemode	1		1	512	1023	0		1	1024	1024	5
uti	write	80 write file	0	la	% osync	0	uniform size dist	1	_	2	1024	2047	0	ç	2	1536	1536	5
trib	mmap write	0 rand write	20	isce	% notification	0	init rate throttle	0	ţi	3	2048	4095	0	Itio	3	2048	2048	5
Dis	rmw	0 append	0	Σ	LRU	0	init read flag	1	ib	4	4096	4096	0	ribr	4	2560	2560	5
u o	mkdir	0 rmdir	0		release version	1			isti	5	4097	8191	0	list	5	3072	3072	5
rati	readdir	0 create	0		Option	Value	Option	Value	e e	6	8192	8192	99	e L	6	3584	3584	5
be	unlink	0 unlink2	0	su.	rand dist behavior	2	% per spot	5	Siz	7	8193	16383	0	· Siz	7	4096	4096	5
e O	stat	0 access	0	cce	min acc per spot	1000	access mult spot	5	fer	8	16384	16384	0	sfei	8	4608	4608	5
Ξ.	rename	0 copyfile	0	A Pa	affinity %	20	spot shape	0	ans	9	16385	32767	0	aŭ	9	5120	5120	5
	locking	0 chmod	0		geometric %	2	align	0	Ě	10	32768	32768	0	ч	10	8192	8192	10
	statfs	0 pathconf	0	S	Option	Value	Option	Value	eac	11	65536	65536	0	Vrit	11	12288	12288	10
ds	Threshold	% Threshold	Value	ter	dedup %	0	dedup within %	0	<u>~</u>	12	98304	98304	0	>	12	16384	16384	10
hol	proc oprate	75 proc latency		Pat	dedup across %	0	dedup group count	1		13	131072	131072	0		13	20480	20480	10
res	global oprate	95 global latanav		ant	dedup granule size	4096	dedup gran rep limit	100		14	262144	262144	0		14	24576	24576	10
Ę	workload variance	5 global latency		nte	compress %	50	comp granule size	8192		15	1048576	1048576	1		15	32768	32768	0
n ers	Parameter	Value Parameter	Value	ပိ	cipher flag	0	pattern version	1										
ete	Procs	1 Dirs per proc	1															

Oprate per proc Avg file size

32 Files per dir

5 200 MiB

EDA_F	RONTEND																	
	Operation	% Operation	%		Option	Value	Option	Value		Slot	Start	End S	%		Slot S	itart	End S	%
	read	0 read file	7		write commit %	15	background	0		0	1	511	4		0	1	511	25
E	mmap read	0 rand read	8		% direct	0	sharemode	0		1	512	1023	0		1	512	1023	10
ntic	write	0 write file	10		% osync	0	uniform size dist	0	_	2	1024	2047	0	_	2	1024	2047	15
rib	mmap write	0 rand write	15	isce	% notification	0	init rate throttle	1	tio	3	2048	4095	2	ţi	3	2048	4095	18
Dist	rmw	0 append	0	Σ	LRU	1	init read flag	0	bu	4	4096	8191	43	ibu	4	4096	8191	27
L L	mkdir	1 rmdir	0		release version	2			stri	5	8192	16383	30	istr	5	8292	16383	3
ratio	readdir	0 create	2	rns	Option	Value	Option	Value	e Di	6	16384	32767	21	e D	6	16384	32767	2
Ibei	unlink	1 unlink2	1	itte	rand dist behavior	0	% per spot	0	Siz	7	32768	65535	0	r Siz	7	32768	65535	0
e	stat	39 access	15	S Po	min acc per spot	0	access mult spot	5	sfer	8	65536	65536	0	sfei	8	65536	65536	0
Ξ	rename	0 copyfile	0	ces	affinity %	0	spot shape	0	rans	9	131072	131072	0	ran	9	131072	131072	0
	locking	0 chmod	1	Ac	geometric %	50	align	0	Гр	10	1	1	0	ē	10	1	1	0
	statfs	0 pathconf	0	ns	Option	Value	Option	Value	Rea	11	1	1	0	Vrit	11	1	1	0
ds	Threshold	% Threshold	Value	ter	dedup %	50	dedup within %	0		12	1	1	0		12	1	1	0
P P	proc oprate	75 proc latency	n/a	Pat	dedup across %	0	dedup group count	1		13	1	1	0		13	1	1	0
res	global oprate	95 global latanav	m / n	ant	dedup granule size	4096	dedup gran rep limit	100		14	1	1	0		14	1	1	0
Ę.	workload variance	5 global latency	n/a	bute	compress %	50	comp granule size	8192		15	1	1	0		15	1	1	0
on ers	Parameter	Value Parameter	Value	ŭ	cipher flag	0	pattern version	2										
utid net	Procs	3 Dirs per proc	10															
xec ran	Oprate per proc	100 Files per dir	10															
Pa E	Avg file size		16 KiB															

EDA_B	ACKEND																	
	Operation	% Operation	%		Option	Value	Option	Value		Slot	Start I	nd	%		Slot 9	Start	End S	%
	read	50 read file	0	SUC	write commit %	15	background	0		0	1	511	0		0	1	511	0
6	mmap read	0 rand read	0	Ĕ	% direct	50	sharemode	0		1	512	1023	0		1	512	1023	0
uti	write	50 write file	0	alla	% osync	5	uniform size dist	0	c	2	1024	2047	0	ç	2	1024	2047	0
trib	mmap write	0 rand write	0	isce	% notification	0	init rate throttle	1	ţi	3	2048	4095	0	Itio	3	2048	4095	0
Dis	rmw	0 append	0	Σ	LRU	1	init read flag	0	libu	4	4096	8191	0	rib	4	4096	8191	0
u o	mkdir	0 rmdir	0		release version	2			list	5	8192	16383	0	Dist	5	8292	16383	0
rati	readdir	0 create	0		Option	Value	Option	Value	e e	6	16384	32767	0	ie D	6	16384	32767	0
be	unlink	0 unlink2	0	ss rns	rand dist behavior	0	% per spot	0	Siz	7	32768	65535	49	r Siz	7	32768	65535	45
<u>e</u>	stat	0 access	0	cce tte	min acc per spot	0	access mult spot	5	sfer	8	65536	65536	51	sfei	8	65536	131072	55
Ē	rename	0 copyfile	0	Pa A	affinity %	0	spot shape	0	ans	9	131072	131072	0	gu	9	131072	131072	0
	locking	0 chmod	0		geometric %	50	align	0	Ę	10	1	1	0	е	10	1	1	0
	statfs	0 pathconf	0	ns	Option	Value	Option	Value	fea	11	1	1	0	Vrit	11	1	1	0
ds	Threshold	% Threshold	Value	ter	dedup %	40	dedup within %	0	<u> </u>	12	1	1	0	>	12	1	1	0
РЧ	proc oprate	75 proc latency	n/a	Pat	dedup across %	0	dedup group count	1		13	1	1	0		13	1	1	0
res	global oprate	95 glabal latanav		at	dedup granule size	4096	dedup gran rep limit	100		14	1	1	0		14	1	1	0
臣	workload variance	5	IId	ht	compress %	20	comp granule size	8192		15	1	1	0		15	1	1	0
ers	Parameter	Value Parameter	Value	ŭ	cipher flag	0	pattern version	2										
utiq	Procs	2 Dirs per proc	5															
xec ran	Oprate per proc	75 Files per dir	10															
Pa E	Avg file size		10 MiB															

SWBUI	LD																	
	Operation	% Operation	%		Option	Value	Option	Value		Slot	Start I	End	%		Slot S	itart	End S	%
	read	0 read file	6		write commit %	33	background	0		0	1	511	1		0	1	511	5
Ę	mmap read	0 rand read	0		% direct	0	sharemode	0		1	512	1023	5		1	512	1023	3
utic	write	0 write file	7		% osync	0	uniform size dist	0		2	1024	2047	7	c	2	1024	2047	10
rib	mmap write	0 rand write	0	isc.	% notification	0	init rate throttle	0	tio	3	2048	4095	7	ţi	3	2048	4095	15
Dist	rmw	0 append	0	Σ	LRU	0	init read flag	1	ibu	4	4096	4096	0	ibu	4	4096	4096	0
l ng	mkdir	1 rmdir	0		release version	1			istr	5	4096	8191	45	istr	5	4096	8191	14
atic	readdir	2 create	1	rns	Option	Value	Option	Value	e D	6	8192	8192	0	e D	6	8192	8192	0
bei	unlink	2 unlink2	0	tte	rand dist behavior	0	% per spot	0	Siz	7	8192	16383	13	Siz	7	8192	16383	7
e O	stat	70 access	6	Ba	min acc per spot	0	access mult spot	5	fer	8	16384	16384	0	sfer	8	16384	16384	0
Ē	rename	0 copyfile	0	ces	affinity %	0	spot shape	0	lans	9	16384	32767	3	ran:	9	16384	32767	6
	locking	0 chmod	5	Ac	geometric %	10	align	0	Fp	10	32768	65535	2	еT	10	32768	65535	4
	statfs	0 pathconf	0	ns	Option	Value	Option	Value	Rea	11	65536	65536	0	Vrit	11	65536	131072	36
ds	Threshold	% Threshold	Value	ter	dedup %	0	dedup within %	0		12	98304	98304	0	-	12	98304	98304	0
loh	proc oprate	75 proc latency	n/a	Pat	dedup across %	0	dedup group count	1		13	65536	131072	17		13	131072	131072	0
Ires	global oprate	95 global latonov	n/2	ent	dedup granule size	4096	dedup gran rep limit	100		14	262144	262144	0		14	262144	262144	0
۲ ۲	workload variance	n/a	ii/a	bit	compress %	80	comp granule size	8192		15	524288	524288	0		15	524288	524288	0
ers	Parameter	Value Parameter	Value	ŭ	cipher flag	0	pattern version	1										
utiq	Procs	5 Dirs per proc	50															
xec	Oprate per proc	100 Files per dir	100															

16 KiB

Par

Avg file size

VDA1																		
	Operation	% Operation	%		Option	Value	Option	Value		Slot	Start	End S	6		Slot	Start	End	%
	read	0 read file	0		write commit %	5	background	0		0	1	511	0		0	1	511	0
E	mmap read	0 rand read	0		% direct	0	sharemode	0		1	512	1023	0		1	512	1023	0
utio	write	100 write file	0		% osync	0	uniform size dist	0	_	2	1024	2047	0	_	2	1024	2047	0
rib	mmap write	0 rand write	0	sce	% notification	0	init rate throttle	0	ig.	3	2048	4095	0	tio	3	2048	4095	0
Dist	rmw	0 append	0	Σ	LRU	0	init read flag	1	ibu	4	4096	4096	0	ibu	4	4096	4096	0
u l	mkdir	0 rmdir	0		release version	1			istr	5	4097	8191	0	istr	5	4097	8191	0
atic	readdir	0 create	0	rns	Option	Value	Option	Value	e	6	8192	8192	0	e D	6	8192	8192	0
ber	unlink	0 unlink2	0	itte	rand dist behavior	0	% per spot	0	Siz	7	8193	16383	0	Siz	7	8193	16383	0
e	stat	0 access	0	Ba	min acc per spot	0	access mult spot	5	sfer	8	16384	16384	0	sfei	8	16384	16384	0
Ξ	rename	0 copyfile	0	ces	affinity %	0	spot shape	0	lans	9	16385	32767	0	ran	9	16385	32767	0
	locking	0 chmod	0	Ac	geometric %	0	align	0	Ψp	10	32768	32768	0	e T	10	32768	32768	5
	statfs	0 pathconf	0	ns	Option	Value	Option	Value	Rea	11	65536	65536	15	Vrit	11	65536	65536	10
ds	Threshold	% Threshold	Value	ter	dedup %	0	dedup within %	0		12	131072	131072	10		12	131072	131072	10
Pol	proc oprate	75 proc latency	n/a	Pat	dedup across %	0	dedup group count	1		13	262144	262144	20		13	262144	262144	25
res	global oprate	95 global latency	n/2	ent	dedup granule size	4096	dedup gran rep limit	100		14	524288	524288	35		14	524288	524288	25
· 두	workload variance	5	11/ d	but	compress %	0	comp granule size	8192		15	1048576	1048576	20		15	1048576	1048576	25
on ers	Parameter	Value Parameter	Value	ŭ	cipher flag	0	pattern version	1										
utic	Procs	1 Dirs per proc	1															
xec Iran	Oprate per proc	9 Files per dir	1															
Pa E	Avg file size		1 GiB															

VDA2																		
	Operation	% Operation	%		Option	Value	Option	Value		Slot	Start	End	%		Slot S	itart	End 🦻	%
	read	5 read file	0		write commit %	0	background	0		0	1	511	0		0	1	511	0
5	mmap read	0 rand read	84		% direct	0	sharemode	0		1	512	1023	0		1	512	1023	0
uti	write	0 write file	0		% osync	0	uniform size dist	0	c	2	1024	2047	0	ç	2	1024	2047	0
trib	mmap write	0 rand write	0	isce	% notification	0	init rate throttle	0	itio	3	2048	4095	0	Ę	3	2048	4095	0
Dis	rmw	2 append	0	Σ	LRU	0	init read flag	1	ibu	4	4096	4096	0	lib	4	4096	4096	0
5	mkdir	0 rmdir	0		release version	1			isti	5	4097	8191	0	list	5	4097	8191	0
rati	readdir	3 create	1		Option	Value	Option	Value	e e	6	8192	8192	0	e e	6	8192	8192	0
be	unlink	1 unlink2	0	ss 'ns	rand dist behavior	0	% per spot	0	Siz	7	8193	16383	0	Siz	7	8193	16383	0
<u>e</u>	stat	2 access	2	tter	min acc per spot	0	access mult spot	5	sfer	8	16384	16384	0	sfei	8	16384	16384	0
Ē	rename	0 copyfile	0	A a	affinity %	0	spot shape	0	aŭ	9	16385	32767	0	an	9	16385	32767	0
	locking	0 chmod	0		geometric %	0	align	0	E E	10	32768	32768	0	e	10	32768	32768	5
	statfs	0 pathconf	0	S	Option	Value	Option	Value	fear	11	65536	65536	15	Vrit	11	65536	65536	10
ds	Threshold	% Threshold	Value	ter	dedup %	0	dedup within %	0		12	131072	131072	10	>	12	131072	131072	10
hol	proc oprate	75 proc latency	n/a	Pat	dedup across %	0	dedup group count	1		13	262144	262144	20		13	262144	262144	25
res	global oprate	95 glabal latanav		ant	dedup granule size	4096	dedup gran rep limit	100		14	524288	524288	35		14	524288	524288	25
f	workload variance	5	li/d	hte	compress %	0	comp granule size	8192		15	1048576	1048576	20		15	1048576	1048576	25
on ers	Parameter	Value Parameter	Value	ŭ	cipher flag	0	pattern version	1										
utiq	Procs	1 Dirs per proc	1															
ran	Oprate per proc	1 Files per dir	1															
Pa E	Avg file size		1 GiB															

VDI																		
	Operation	% Operation	%		Option	Value	Option	Value		Slot	Start	End	%		Slot	Start	End	%
	read	6 read file	0		write commit %	100	background	0		0	512	512	1		0	512	512	21
Ę	mmap read	0 rand read	20		% direct	100	sharemode	0		1	2048	2048	1		1	1024	1024	2
utic	write	9 write file	0		% osync	0	uniform size dist	0		2	2560	3584	1	e	2	2048	2048	1
rib	mmap write	0 rand write	64	isce	% notification	0	init rate throttle	0	tio	3	4096	4096	20	tion	3	4096	4096	47
Dist	rmw	0 append	0	Σ	LRU	0	init read flag	1	ibu	4	4608	7680	1	ibu	4	8192	8192	6
u n	mkdir	0 rmdir	0		release version	1			istr	5	8192	8192	4	istr	5	8704	15360	3
atio	readdir	0 create	0	rns	Option	Value	Option	Value	e	6	8704	15872	4	e D	6	16384	16384	5
bei	unlink	0 unlink2	0	itte	rand dist behavior	0	% per spot	0	Siz	7	16384	16384	42	Siz	7	16896	30720	5
<u> </u>	stat	0 access	1	Ba	min acc per spot	0	access mult spot	5	sfer	8	16896	32256	3	sfei	8	32768	32768	1
Ξ	rename	0 copyfile	0	ces	affinity %	0	spot shape	0	rans	9	32768	32768	14	ran	9	35328	64000	3
	locking	0 chmod	0	Ac	geometric %	90	align	0	Р	10	33280	64024	1	Ē	10	65536	65536	2
	statfs	0 pathconf	0	ns	Option	Value	Option	Value	Rea	11	65536	65536	6	Nrit	11	69632	126976	1
ds	Threshold	% Threshold	Value	ter	dedup %	0	dedup within %	0		12	66048	126976	1	-	12	131072	131072	3
loh	proc oprate	75 proc latency	n/a	Pat	dedup across %	0	dedup group count	1		13	131072	131072	1		13	262144	262144	0
Ires	global oprate	90 global latonov	n/2	ent	dedup granule size	4096	dedup gran rep limit	100		14	262144	262144	0		14	524288	524288	0
Ē	workload variance	n/a global latency	ny a	bit	compress %	60	comp granule size	8192		15	524288	524288	0		15	1048576	1048576	0
on ers	Parameter	Value Parameter	Value	ŭ	cipher flag	0	pattern version	1										
utiq	Procs	2 Dirs per proc	1															
ian ran	Oprate per proc	100 Files per dir	1															

Avg file size

500 MiB

WORK	OAD NAME																	
	Operation	%	Operation	9	6	Option	Value	Option	Value		Slot	Start	End	%		Slot Start	End	%
	read		read file		Sn	write commit %		background			C)				0		
Ę	mmap read		rand read			% direct		sharemode			1	-				1		
utio	write		write file		a la	% osync		uniform size dist		_	2	2			_	2		
rib	mmap write		rand write		SCe	% notification		init rate throttle		ţi	3	5			tio	3		
Dist	rmw		append		Σ	LRU		init read flag		- Ing	4	Ļ			ibu	4		
u I	mkdir		rmdir			release version				stri	5	i			istr	5		
atic	readdir		create		rns	Option	Value	Option	Value	e	e	5			e D	6		
bei	unlink		unlink2		itte	rand dist behavior		% per spot		Siz	7	,			, Siz	7		
<u>e</u>	stat		access		Pa Ba	min acc per spot		access mult spot		sfer	8	3			sfei	8		
Ē	rename		copyfile		ces	affinity %		spot shape		rans	ç)			ran	9		
	locking		chmod		Ac	geometric %		align		Г Р	10)			еT	10		
	statfs		pathconf		ns	Option	Value	Option	Value	Rea	11	-			Vrit	11		
ds	Threshold	%	Threshold	Valu	e j	dedup %		dedup within %			12	2			-	12		
Pol	proc oprate		proc latency		Pat	dedup across %		dedup group count			13	;				13		
ires	global oprate		global latoncy		ent	dedup granule size		dedup gran rep limit			14	Ļ				14		
Ē	workload variance		giobal latelicy		Ť	compress %		comp granule size			15	5				15		
on ers	Parameter	Value	Parameter	Valu	e Ö	cipher flag		pattern version										
let	Procs		Dirs per proc															
an	Oprate per proc		Files per dir															

Avg file size